## University of Alberta

## Collusion Detection in Sequential Games

by

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To my dear family, for being there.

## Abstract

Collusion is the deliberate cooperation of two or more parties to the detriment of others. While this behaviour can be highly profitable for colluders (for example, in auctions and online games), it is considered illegal and unfair in many sequential decision-making domains and presents many challenging problems in these systems.

In this thesis we present an automatic collusion detection method for extensive form games. This method uses a novel object, called a collusion table, that aims to capture the effects of collusive behaviour on the utility of players without committing to any particular pattern of behaviour. We also introduce a general method for developing collusive agents which was necessary to create a dataset of labelled colluding and normal agents. The effectiveness of our collusion detection method is demonstrated experimentally. Our results show that this method provides promising accuracy, detecting collusion by both strong and weak agents.

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## Chapter 1

## Introduction

Collusion is the practice of two or more parties deliberately cooperating to the detriment of other parties. While this cooperation is allowed and encouraged in some multi-agent settings, in other domains it is illegal or prohibited and can provide colluders an unfair advantage. Because of this, detecting and preventing collusion is a challenge of major interest in many different settings.

One real world system in which collusion presents a challenging problem is financial markets. Collusive behaviour in these domains includes insider trading and market manipulation to set prices or limit production. Previous research has introduced some collusion detection methods that utilize pre-established parameters in computerized trading systems as an indicator of normal versus suspicious activities. If these parameters' normal values are exceeded in any trade, the system identifies and reports the trade [21]. This mechanism is based on parameters which are domain specific and thus cannot be generalized to other settings where collusion detection is desired. Additionally, the system does not completely determine that illegal activity has occurred, but further consideration of trained staff is needed to fully detect that unlawful interaction occurred.

Another setting in which collusive behaviour is illegal is games such as poker. On online poker websites participants are instructed to report any suspicious behaviour to the website administrators. Additionally, some detection methods are used which warns security personnel if any pre-specified unusual play patterns occur [15]. These methods are also based on known patterns of collusion and require human experts to examine the history of the games in which suspected colluders
participated. When the number of players increases or the volume of games is large, in-depth examination of data for collusion patterns costs an enormous amount of resources, both in terms of human experts' time and money.

Some detection methods build a model of normal or collusive behaviour based on previous interactions of agents [25,26]. However, such models can only detect collusive behaviour that matches the known collusive model or differs from the normal model and cannot identify other collusive behaviours.

While no collusion detection system can completely replace human experts, automated detection methods can assist them and improve the efficiency of investigations. Ideally, an automated collusion detection system should be able to detect collusive behaviour based not only on a specific pattern but also be able to detect unknown patterns that are also collusive. The method should depend on the actions taken by agents and not on any domain specific factors. The goal of this thesis is to design an automated general-purpose collusion detection method that can be used in different multi-agent systems.

Note that we distinguish between the concept of cheating and colluding. We define cheating as using information or actions that is not permitted by the structure and rules of the game. For example, when players share their private cards using a back channel, it is cheating. However, collusion is a series of valid actions which leads to an unfair advantage for colluders. For example, when a player reraises his partner to increase the leverage of his bet, it is colluding. The focus of this work is collusion detection, although we hypothesize that our technique could detect cheating as well.

Our approach to developing a general-purpose collusion detection method focuses fundamentally on our definition of collusion: that colluders act to increase their joint utility. First, we define the collusion detection problem in the domain of extensive form games, a general model of sequential multi-agent decision making systems. We then focus on how the actions taken by players affect the utility of all participants. This information is summarized in a novel data structure called a collusion table. We then evaluate our method by a series of experiments in the domain of poker, where collusive behaviour can give colluders a great advantage and col-
lusion detection is a real world challenge. The experiments validate our approach for detecting collusion in the domain of poker. This method can also be applied to other domains since it does not use any domain specific knowledge.

This thesis will proceed as follows. First, in Chapter 2, we describe work related to collusion detection, which can be divided into two categories: the first is previous work based on known patterns of collusion or domain specific methods. The second is work in the domain of operations research and game theory for evaluating agent behaviour. In 2006, Zinkevich and colleagues proposed advantage sum estimators for constructing low-variance unbiased estimates of an agent's performance [27]. This approach has been utilized in agent evaluation techniques like DIVAT and MIVAT which will be explained in Chapter 2. Unlike other methods of estimating agent's performance, this method examines the whole history of players' interactions. We also employ the same underlying idea as a basis for our method of evaluating the effect of agents' actions on each other in games.

Chapter 4 is devoted to the construction of collusion tables. First, the design and semantics of collusion tables are described, followed by some examples of collusion tables. We also introduce a number of different scoring methods which are designed to evaluate the likeliness that a pair of agents are colluding, given a collusion table. While collusion tables could be constructed in many different ways, in this work we use value functions from histories of the game to real numbers. This approach has been shown to be successful in the field of agent evaluation [2,23,27]. We propose three techniques for constructing a value function in Chapter 5.

Ideally, a dataset of play by both colluding and non-colluding agents, where each agent is labelled as such, would be used to validate our approach. It would be most desirable for this dataset to be constructed from human play. However, as such a dataset is not available, we instead created a synthetic dataset of matches using various "bots". This requires developing poker bots that collude with partners, which we did by modifying the utility function of the colluders so that they consider their partner's utility as their own. This technique is described in Chapter 3 and is accompanied by the results of experiments that verify the effectiveness of our method at creating profitable colluding bots. The different kinds of agents which
we developed to represent players with different skill levels and a description of our experimental design is described in Chapter 6. The results of our experiments are presented in Chapter 7. Finally, we conclude this work in Chapter 8 and discuss possible future work.

## Chapter 2

## Background and Related Work

In this chapter, we define the general domain in which collusion detection is investigated. Then, the features of poker which make it a proper experimental domain are explained. After presenting the problem formulation, previous work on collusion detection and other related work in the field of game theory is described. Finally, a state of the art technique for developing poker agents is explained which will be used in our experiments.

### 2.1 Definitions and Notation

In this section the preliminary definitions used in this thesis are presented.

### 2.1.1 Extensive form games

Definition 1. (Adapted from Osborne and Rubinstein [14] ) A zero-sum Extensive form game is defined to have the following components:

- A finite set $N$ (the set of players).
- A finite set $H$ of sequences such that : (1) the empty sequence $\emptyset$ is a member of $H$. (2) If $h^{\prime}$ is a prefix of $h$ and $h \in H$, then $h^{\prime}$ is also a member of $H$. (3) If every subsequence $h^{\prime}$ of an infinite sequence $h$ is a member of $H$, then $h \in H$.

Each member of $H$ is called a history. $Z \subseteq H$ is the set of terminal histories which are sequences that are infinite or are not a subsequence of any other sequence in $H$.

- For each player $i \in N$ a function $u_{i}: Z \mapsto \mathbb{R}$ which assigns a utility for player $i$ to each terminal history $z \in Z$.
- A finite set $A$ consisting of all possible actions a player can take. $A(h)=$ $\{a \in A: h a \in H\}$ denotes the set of actions ${ }^{1}$ that a player can take after $a$ history $h \in H$.
- A player function $P: H \backslash Z \mapsto N \cup\{c\}$ which assigns a player (from set $\{1,2, \ldots, N\}$ or the chance player c) to each non-terminal history $h . P(h)$ is the player who takes an action after history $h$. If $P(h)=c$ then chance determines the action taken after the non-terminal history $h$.
- A function $\sigma_{c}$ which associates with every history in $\{h \in H: P(h)=c\}$ a probability distribution $\sigma_{c}(. \mid h)$ on $A(h) . \sigma_{c}(a \mid h)$ is the probability that action a occurs after the history $h$.
- For each player $i \in N$ a partition $\mathcal{I}_{i}$ of $\{h \in H: P(h)=i\}$ such that $A(h)=A\left(h^{\prime}\right)$ whenever $h$ and $h^{\prime}$ are in the same member of the partition. $\mathcal{I}_{i}$ is the information partition of player $i$; a set $I_{i} \in \mathcal{I}_{i}$ is an information set of player $i$.

A game is called a zero-sum game if we have $\sum_{i \in N} u_{i}=0$ for all terminal histories in $Z$. Extensive form games can be represented as a game tree. Each node represents a player (or chance player) and each edge represents a valid action. The path from root to each node determines a history of the game. All possible sequences of actions are represented in a game tree (terminal histories with their corresponding utilities). In extensive form games with perfect information (like chess), a player can determine the state of the game and has complete information about all players' actions taken previously (by looking at the board). However, in extensive form games with imperfect information, a player may not have information about previous actions that are taken by players or chance. This means that players may not be able to differentiate between game states. Consider the example shown in Figure

[^0]2.1. In this game tree, player 1 is the first player to act and can take action $A$ or action $B$. Then player 2 is the player to act who can choose action $C$ or action $D$. There are five terminal histories in this game which are marked as black. In this figure the dashed line indicates that histories AC and AD are in the same information set of player 1 which means this player cannot distinguish between history AC and history AD. In either case, it is player 1's turn to play and choose an action from set $\{E, F\}$.


Figure 2.1: An example of an extensive game with imperfect information.

### 2.1.2 Poker and its properties

In this thesis, poker is used as an experimental domain to validate our collusion detection methods. Some of the properties that make poker a suitable domain for this research are as follows:

- Poker is a multiplayer game with imperfect information. Important information which a player cannot access is the private cards of opponents. Therefore, a player cannot make decisions based on this information unless he illegally accesses it by cheating or colluding.
- Poker is a popular game; not only because it is entertaining, but also because money is involved in poker. Many poker experts view the game as a way to gain money. This means that players may be more motivated to collude or cheat in poker than other games.
- Poker players have many different skill levels. Also, there is no single, known winning strategy in this game. This results in a wide variety of poker playing styles which complicate the task of detecting unusual behaviours which might be signs of collusion or cheating.
- Finally, poker provides us a good representation of real world environments. Settings like auctions and market places are very similar to poker in their structure. Therefore, techniques that we develop for poker should be extendable to these real-world settings as well.


### 2.1.3 Strategy and Nash Equilibrium

The outcome of an extensive form game is highly dependent on the strategies of the players. Informally, the actions that a player chooses at different states of a game is called that player's strategy. Each strategy is simply a set of action distributions for each information set in the game. Since a player's strategy affects the outcome of the game, a player would like to choose a strategy which increases the player's utility at reached outcomes.

Definition 2. (Adapted from Osborne and Rubinstein [14] ) A behavioural strategy of player $i$ is a collection $\left(\beta_{i}\left(I_{i}\right)\right)_{I_{i} \in \mathcal{I}_{i}}$ of independent probability measures, where $\beta_{i}\left(I_{i}\right)$ is a probability measure over $A\left(I_{i}\right)$.
$\beta$ indicates a profile of behavioural strategies such that for each player $i, \beta_{i} \in \beta$ is the behavioural strategy of player $i . \beta_{-i}$ indicates a strategy profile for all the players in a game except for player $i$. The expected utility of player $i$ if all the players in the game follow the strategy profile $\beta$ is $u_{i}(\beta)$.

Definition 3. (Adapted from Osborne and Rubinstein [14] ) A Nash equilibrium in behavioural strategies of an extensive game is a strategy profile $\beta^{*}$ of behavioural strategies with the property that for every player $i \in N$ we have

$$
u_{i}\left(\beta_{-i}^{*}, \beta_{i}^{*}\right) \succsim_{i} u_{i}\left(\beta_{-i}^{*}, \beta_{i}\right) \text { for every behavioural strategy } \beta_{i} \text { of player } i .
$$

In this thesis, we refer to a behavioural strategy as a strategy. An entity with a strategy to play as a player or position in the game is called positional agent or player. A set of positional agents which has at least one strategy for each position in the game forms an agent. Thus, an agent can participate as any player in a game.

### 2.1.4 The Collusion detection problem

In this section we introduce the problem of collusion detection. Assume that we have a dataset of past interactions between agents from some population of agents $M$.

Definition 4. A game episode $g$ is a tuple $\left\langle P_{g}, \phi_{g}, z_{g}\right\rangle$, in which $P_{g} \subseteq M$ is the subset of agents participating in $g$. A mapping function $\phi_{g}: P_{g} \mapsto N$ that associates each agent with a player or position in the game and $z_{g}$ is the terminal history reached at the end of game episode $g$.

Definition 5. Given a set of agents M participating in a dataset of game episodes $D$, the Collusion Detection Problem is to find a function $\vartheta: M^{2} \mapsto \mathbb{R}$ that measures the possibility of collusion for each pair of agents in $M$ and ranks them accordingly.

That means, if $\vartheta(a, b)>\vartheta(c, d)$, then it is more likely that agents $(a, b)$ are colluding than agents $(c, d)$ are. In other words, the goal of collusion detection is to rank the pair of players participating in a dataset in order of their likeliness of collusion.

### 2.2 Experimental Domain

As described in section 2.1.2, poker has properties which make it a suitable experimental domain for the purpose of this research. In this section, we explain the variation of poker that is used in our experiments.

### 2.2.1 [2-4] Limit Hold'em Poker

[2-4] Hold'em Poker is a small version of Texas Hold'em Poker. In Texas Hold'em, we have 2 to 10 players. Each player is given a hand of 2 private cards from a
shuffled deck of 52 cards. There are four rounds in the game, pre-flop, flop, turn and river. A specific number of public cards are revealed after each round (three after the pre-flop, one after the flop and one after the turn). After each deal, players have the options to bet (add more money to the pot), check/call (match the money that other players put in the pot) or fold (discard their hand and lose the money they already put in the pot). The objective is to win the money in the pot. After the last round, player who has not folded that can make the best 5-card poker hand out of his private cards and the public cards wins the pot.

Texas Hold'em is usually played using small and big blind bets. These are the first two bets which the first two players must begin with. A dealer button is used to specify the position of the dealer. The player to the left of dealer bets the small blind and the big blind is bet by the player to the left of the small blind. The dealer button rotates clockwise after each game episode to change the position of the dealer and blinds. In this way, the effect of position on the utility of the players is averaged out over the hands.

There are two variants of Texas Hold'em Poker and [2-4] Hold'em Poker called limit and no-limit. In the limit version, four raises are allowed in each round of the game. Additionally, the size of the bets is fixed. On the other hand, in each round of no-limit games, players can raise any number of times by any arbitrary amount which is greater than or equal to the minimum bet size and less than their total number of chips. Therefore, the size of the game tree in no-limit games is much larger than in limit games.
[2-4] Hold'em Poker is exactly like limit Texas Hold'em except for the number of rounds. The game ends after the flop betting round, and there are no turn or river rounds. Therefore, the number of public cards in total is three (The name, [2-4], refers to two rounds and four raises that are allowed in each round). The game that is used in our experiments is [2-4] Limit Hold'em poker because of its similarity to Texas Hold'em and its smaller size. The size of the small and big blinds are set to 5 and 10 , respectively.

### 2.3 Previous Work

There have been a few previous studies on detecting and preventing collusion in different domains. Many suggested methods are based on patterns of collusion which are specified by human experts. The others design methods based on features of the game which are mostly domain specific. In this section an overview of this work in the domain of auctions and games is presented .

Auctions are one of the important areas in which collusion is prohibited and thus has received great attention by scientists. Most of this work, however, discusses the problem of collusion prevention and is aimed at designing auctions so that collusion is not a profitable strategy for bidders [7, 10, 13]. In 1989, Hendricks and Porter [8] argued that the presence of collusive behaviour is highly dependent on the object being auctioned and the auction rules.

Robinson [17] compared the stability of cartels in oral auctions versus sealed high-bid auctions. The results shows that cartels are stable in oral auctions but not in sealed high-bid auctions, but this doesn't mean that collusion can never happen in the latter. Bachrach et al. [1] investigated collusive behaviour of bidders in a class of auctions called Vickrey-Clarke-Groves auctions using a cooperative game theory approach. They examined different auctions with respect to the possibility of a stable agreement between colluders. They showed that in some auctions participants may not be able to form a long-lasting agreement even if they have complete information about each other's preferences.

Another area in which collusion is a real-world challenge and has been studied is online games. In 2010, Yan introduced the problem of detecting collusion in online bridge as a hard problem due to anonymity and the benefit of sharing information through back channels that players have when playing over internet [26]. The motivation was to solve the general problem of detecting the usage of prohibited information in decision making, which is categorized in our termination as cheating not colluding. The approach suggested by Yan is to focus on critical parts of the play in which traces of cheating may appear more and compare the sequence of decisions of the players with a model of non-colluding play. However, their ap-
proach is completely based on the "critical parts" they introduced in bridge and is not applicable to other domains.

Smed et al. [19,20] also investigated the problem of collusion in different games and gave an extensive classification of collusion based on types of agreement that colluders can have. They also introduced a simplified version of the game of PacMan for evaluating collusion detection methods. Laasonen et al. [11] utilized this game for their experiments to detect features that are indications of collusion. The features they suggested are secondary factors which are domain specific. Finally, they gave an analysis of the utility functions of different groups of colluders in the designed experimental settings.

### 2.4 Related Work on Agent Performance Evaluation

Performance evaluation is a well-studied problem which has some similarities to the problem of detecting collusion. This problem arises in any multi-agent system in which designing better agents is needed. Because of the stochastic nature of the environment and agent's decisions, designing a low variance estimator of agent performance has been the target of several recent studies. These issues also arise in detecting collusion and makes the problem of designing a collusion detector method a challenge. In 2006, Zinkevich et al. [27] introduced a low variance estimator called an advantage sum. The principle of this method has been used in two other agent evaluation techniques on which the idea of our collusion detection method is built. In this section we briefly explain these methods (for a complete description of these methods, see the original papers [2, 23, 27]).

### 2.4.1 Advantage Sum Estimators and DIVAT

In 2006, Zinkevich et al. [27] showed that given any value function on the histories of a finite extensive form game, an unbiased estimator can be generated. They do this as follows: Assume one is given a real valued function on histories $V_{i}: H \mapsto \mathbb{R}$ such that $V_{i}(z)=u_{i}(z) \quad \forall z \in Z$ for any given player $i \in N$. Define the following real-valued functions on terminal histories:

$$
\begin{gather*}
S_{V_{i}}(z)=\sum_{\substack{h a \sqsubseteq z \\
P(h) \neq c}} V_{i}(h a)-V_{i}(h)  \tag{2.1}\\
L_{V_{i}}(z)=\sum_{\substack{h a \sqsubseteq z \\
P(h)=c}} V_{i}(h a)-V_{i}(h)  \tag{2.2}\\
\operatorname{Pos}_{V_{i}}=V_{i}(\emptyset) \tag{2.3}
\end{gather*}
$$

We write $h a \sqsubseteq z$ to denote that history $h a$ is a prefix of terminal history $z$. We call these values skill, luck and position of player $i$ (note that in this definition skill is computed with respect to the players participating in the game and will change if the players change). The utility can be rewritten as a function of skill, luck and position because terms cancel when summing skill, luck and position.

$$
\begin{array}{rc}
S_{V_{i}}(z)+L_{V_{i}}(z)+\operatorname{Pos}_{V_{i}}= & \\
& V_{i}(\emptyset) \\
& +V_{i}\left(a_{1}\right)-V_{i}(\emptyset) \\
& +V_{i}\left(a_{1} a_{2}\right)-V_{i}\left(a_{1}\right) \\
+\cdots+ & V_{i}(z)-V_{i}\left(a_{1} \cdots a_{|z|-1}\right) \\
& =V_{i}(z) \\
& =u_{i}(z) \tag{2.4}
\end{array}
$$

The advantage sum estimator is then defined as follows:

$$
\begin{equation*}
\hat{u}_{V_{i}}(z)=S_{V_{i}}(z)+\operatorname{Pos}_{V_{i}} \tag{2.5}
\end{equation*}
$$

To have an unbiased estimator, Zinkevich et al. [27] suggest choosing the value function for player $i$ so that the expected value of luck for player $i$ equals zero. This is called the zero-luck constraint. In this case, we have:

$$
\begin{array}{rlr}
E\left[\hat{u}_{V_{i}}(z) \mid \sigma\right] & =E\left[S_{V_{i}}(z)+\operatorname{Pos}_{V_{i}} \mid \sigma\right] \\
& =E\left[S_{V_{i}}(z)+L_{V_{i}}(z)+\operatorname{Pos}_{V_{i}} \mid \sigma\right] \quad \text { if } E\left[L_{V_{i}}(z) \mid \sigma\right]=0 \\
& =E\left[u_{V_{i}}(z) \mid \sigma\right] &
\end{array}
$$

Specifically, they suggested that if the value function is chosen such that the value of histories before a chance node is equal to the expected value of the histories right after a chance node for player $i$, then the zero-luck constraint is satisfied. This approach will be used later in the computation of collusion values in Chapter 4.

Based on this technique, Billings and Kan proposed the Ignorant Value Analysis Tool (DIVAT), a low variance estimator of agents in two player Texas Hold'em poker which uses a hand designed value function which satisfies the zero-luck constraint [2].

### 2.4.2 MIVAT

Although DIVAT has been shown to have very good performance in two player Texas Hold'em limit poker, it cannot be expanded to other domains since it uses a hand crafted value function. White and Bowling proposed a more general approach to utilize the advantage sum technique called the Informed Value Assessment Tool. Instead of using a domain specific hand designed value function, they learn a value function based on the features of the domain, given sample data of previous matches played by players, and so is informed by data of past interactions [23].

To satisfy the zero-luck constraint, they assume that the value function is only learned for histories following a chance node. Then, the value of the game for player $i$ at histories where chance is next to act is defined to be computed from a weighted sum of the other histories. The value function is then:

$$
V_{i}= \begin{cases}\sum_{a \in A(h)} \sigma_{c}(a \mid h) V_{i}(h a) & \text { if } P(h)=c, \\ \text { learned function } & \text { Otherwise }\end{cases}
$$

Given $T$ samples of outcomes, they minimized the sum of the estimated utility variance over $T$ samples. To make this optimization tractable, they focused on
the class of linear value functions and proposed a closed form solution for such value functions. Our method extends these approaches to the problem of collusion detection.

### 2.5 Counterfactual Regret Minimization

Regret minimization is a well-known concept in online learning [3]. In 2008, Zinkevich et al. [28] proposed a technique for solving extensive form games based on regret minimization. They introduced a new concept called counterfactual regret and proved that minimizing counterfactual regret leads to minimizing overall regret which results in an approximate Nash equilibrium strategy in two-player zero-sum games. This approach also requires only memory linear in the number of information sets instead of game states. As the agents used in our experiments are built using this technique, we briefly describe it in this section. ${ }^{2}$

The concept of regret is similar to opportunity cost in economics. We can informally define it as follows: suppose one has taken action $a$ and gained utility $u(a)$. Regret is the different between the utility that you gained and the maximum utility that could have been possible for you to gain if you had taken the right action $a^{*}$ or $u\left(a^{*}\right)-u(a)$. A strategy selection algorithm is called regret minimizing if the average overall regret of player $i$ playing the chosen strategy $\sigma_{i}^{t}$ in iteration $t$ goes to zero as the number of repeated choices played goes to infinity.

Counterfactual Regret Minimization is a regret minimizing algorithm for solving zero-sum extensive form games. In this algorithm, positional strategies play repeated games against each other. The algorithm's action probability distribution is initiated uniformly from all possible choices in all information sets. In each round and for each information set, a positional agent improves its strategy (by changing the probability distribution over actions) so as to minimize the regret of its subtree given that strategies of other participants are fixed. In a two player zero-sum game, it is shown that in self-play as the number of games increases, the regret

[^1]minimizing behaviour of positional strategies will cause them to approach a Nash equilibrium [9]. CFR-generated strategies have been shown to have good performance in multiplayer poker games (e.g. it achieved first place in the three-player limit Texas Hold'em category at the Annual Computer Poker Competition in 2012 ${ }^{1}$ [6]). We used CFR in our experiments due to having the advantage of its low memory requirement and possibility of parallel computation [9] which improves the speed of the program.

[^2]
## Chapter 3

## Colluding Bots

In order to evaluate our proposed collusion detection methods, a synthetic dataset of labeled colluding and non-colluding agents is required. Since a labeled human dataset of poker games is not available, we used a population of "labeled" bots to play a set of game episodes. While there have been a number of non-colluding bots developed by the CPRG and other research groups [4], to our knowledge no colluding agent has ever been developed. Thus, to create a synthetic dataset, we first develop colluding bots.

There are many opportunities in most multi-agent settings for players to collude. Some of these collusion methods depend on the system's characteristics and cannot be utilized in other domains. Here, we propose a collusion method which is not dependent on any system properties or constrained to any particular pattern of collusion. Our method centers on the very definition of collusion: that colluding is jointly beneficial for all colluding partners.

After describing our method we evaluate it in the domain of [2-4] Hold'em poker and the results show that agents are indeed successfully colluding. However, this method could be used in any multi-agent system, including any other variations of poker. In this chapter our technique for developing colluding bots is introduced and the experimental results which validates the method are presented.

### 3.1 Learning to Collude: Modified Utility Function

In a multi-agent system players can collude using various methods. Many of these techniques are domain dependent. That is, colluders take advantage of characteristics of the game to increase their utilities. Other methods utilize the environment to collude. For instance, using an independent communication channel to share private information is one method of colluding in online poker. However, all of these methods have one thing in common: the purpose is to increase the joint utility of the colluders. One way this can be interpreted is that each party involved in collusion partially considers his partner's utility as his own and plays accordingly. For instance, consider two colluding players in poker, one with a weak and one with a strong hand. One example of collusive behaviour would be to gain the pot by raising and re-raising until all non-colluding players with marginal hands fold. After that, the colluder with the weak hand folds to a raise and the colluder with the strong hand gets the pot. This method is known as active collusion in which the weak colluder considers the utility of the strong colluder as his own completely [24]. In contrast, if he were to play normally, the best strategy would be to fold at the beginning of the game to avoid a huge loss.

We use a utility function to capture the notion of colluding agents (partially) considering their partner's utility as their own. This is done by modifying the utility function. Once the utility function is altered, we can use any strategy creation method to build the strategies for each position in the modified game. We define the modified utility function for colluding player $i$ who colludes with player $j$ as:

$$
\hat{u}_{i}(z)=u_{i}(z)+\lambda u_{j}(z) \quad \text { for all } z \in Z,
$$

where parameter $\lambda$ specifies how much a player considers his partner's utility as his own. We then apply the Counterfactual Regret Minimization (CFR) method [16,28] to the modified game to create positional collusive strategies. CFR is used since it has been shown to generate winning computer poker agents for multiplayer games [16]. For each combination of positions in the game, a pair of colluding positional strategies were created. Due to the huge number of game states in poker, we first abstract the game and then develop colluding strategies in the abstracted game. In
general the less abstraction is used, the higher the quality of the created strategy will be in the full game (although some examples of abstraction pathologies in toy poker games exists [22]).

### 3.2 Results and Discussion

To determine which value of $\lambda$ was most beneficial for the colluders, we ran several sets of experiments. First, positional strategies for different $\lambda$ values were created for the game of [2-4] Hold'em poker. To ensue the consistency of the developed positional strategies, different random seeds were used when building the strategies using CFR. Once the strategies were created, a set of experiments was ran to verify that players can collude. Each experiment consists of three players, two colluders and one normal player who played one million hands (game episodes) of [2-4] Hold'em poker. The dealer button did not rotate in these game sets since the players are position dependent. Table 3.1 shows the results of this set of games: sixteen matches, each of one million hands. For each value of $\lambda>0$, a one million hand match was played for each of the three different combinations of colluders' and non-colluder's positions. The performance of players is measured in milli-bigblind/game ( $\mathrm{mbb} / \mathrm{g}$ ), where milli-big-blind is 0.001 big blind. The name of the players in this table shows the player index followed by its partner in that match. When a player does not have a partner, $N$ is used to show that it is a normal noncolluding player. For instance, player 1,2 on the first row of the table indicates that player 1 is colluding with player 2 in this match.

We also ran a one-million-hand match in which nobody is colluding (that is when $\lambda=0$ ) for the sake of comparison. In this set rotating positional strategies does not provide further information. The results for this match is presented in the last three rows of Table 3.1.

Table 3.1 demonstrates the average winnings of each player during a one-millionhand set. As described in Section 2.4, the utility of a player can be divided into the utility achieved from skill of the player, luck, and the benefit of position. Table 3.1 shows that the player in position three (the last player to act) always has a higher

| $\lambda$ | Player | Mean (mbb/g) | Sample Standard Deviation | *Estimated SD 95\% CI |
| :---: | :---: | :---: | :---: | :---: |
| 0.3 | 1,2 | -93.7215 | 2711.04 | 5.31364 |
|  | 2,1 | -177.131 | 2974.27 | 5.82958 |
|  | 3,N | 270.852 | 2487.06 | 4.87464 |
|  | 1,3 | -92.152 | 2742.66 | 5.37561 |
|  | 2,N | -186.153 | 2960.24 | 5.80207 |
|  | 3,1 | 278.305 | 2563.64 | 5.02473 |
|  | 1,N | -96.8015 | 2709.56 | 5.31075 |
|  | 2,3 | -179.526 | 2945.62 | 5.77342 |
|  | 3,2 | 276.327 | 2456.56 | 4.81486 |
| 0.5 | 1,2 | -93.2045 | 2723.4 | 5.33787 |
|  | 2,1 | -174.442 | 2989.94 | 5.86029 |
|  | 3,N | 267.646 | 2484.89 | 4.87039 |
|  | 1,3 | -86.3735 | 2779.6 | 5.44801 |
|  | 2,N | -192.334 | 2976.09 | 5.83313 |
|  | 3,1 | 278.708 | 2644.37 | 5.18296 |
|  | 1,N | -104.814 | 2707.04 | 5.3058 |
|  | 2,3 | -170.766 | 2912.22 | 5.70796 |
|  | 3,2 | 275.579 | 2386.75 | 4.67802 |
| 0.9 | 1,2 | -83.2175 | 2823.85 | 5.53474 |
|  | 2,1 | -176.58 | 3092.45 | 6.0612 |
|  | 3,N | 259.798 | 2510.98 | 4.92153 |
|  | 1,3 | -298.743 | 4026.31 | 7.89157 |
|  | 2,N | -282.812 | 2995.33 | 5.87085 |
|  | 3,1 | 581.554 | 3793.39 | 7.43504 |
|  | 1,N | -153.058 | 2794.93 | 5.47806 |
|  | 2,3 | -88.525 | 2981.18 | 5.84311 |
|  | 3,2 | 241.583 | 2448.42 | 4.79891 |
| 0.99 | 1,2 | -69.4655 | 2901.35 | 5.68665 |
|  | 2,1 | -186.345 | 3153.23 | 6.18032 |
|  | 3,N | 255.811 | 2528.09 | 4.95505 |
|  | 1,3 | -182.321 | 3728.09 | 7.30706 |
|  | 2,N | -259.239 | 3057.01 | 5.99173 |
|  | 3,1 | 441.561 | 3639.56 | 7.13355 |
|  | 1,N | -155.818 | 2812.48 | 5.51246 |
|  | 2,3 | -72.5558 | 3077.89 | 6.03267 |
|  | 3,2 | 228.374 | 2576.63 | 5.0501 |
| 0.999 | 1,2 | -67.7375 | 2908.19 | 5.70005 |
|  | 2,1 | -187.542 | 3159.99 | 6.19357 |
|  | 3,N | 255.28 | 2530.73 | 4.96024 |
|  | 1,3 | -175.559 | 3706.38 | 7.26451 |
|  | 2,N | -260.049 | 3060.46 | 5.9985 |
|  | 3,1 | 435.608 | 3626.23 | 7.1074 |
|  | 1,N | -153.228 | 2815.2 | 5.5178 |
|  | 2,3 | -71.6845 | 3079.1 | 6.03503 |
|  | 3,2 | 224.913 | 2573.54 | 5.04414 |
| 0 | 1,N | -94.722 | 2707.84 | 5.30737 |
|  | 2,N | -181.684 | 2956.63 | 5.79499 |
|  | 3,N | 276.406 | 2491.3 | 4.88295 |

Table 3.1: Performance of Players for a set of matches with different lambda values. (* Estimated Standard Deviation for a 95\% Confidence Interval.)
value in comparison to positions one and two. This reveals the effect of position on the utility of a player. Remarkably, the player in position three wins the most even when the other two players are colluding against him. However, collusion helps players in positions one and two to lose less than when they are not colluding. A summary of the colluders' average winnings when in the different seats is presented in Table 3.2. These results reveals that the most beneficial collusion occurs when players are in positions one and three.

| $\lambda$ | Colluders' Average Wining for Positions: |  |  | Average Colluders' <br> Wining with Rotation |
| :---: | :---: | :---: | :---: | :---: |
|  | $1 \& 2$ | $1 \& 3$ | $2 \& 3$ | 0.00000 |
| 0 | -276.406 | 181.684 | 94.722 | 4.03383 |
| 0.3 | -270.8525 | 186.153 | 96.801 | 9.83367 |
| 0.5 | -267.6465 | 192.3345 | 104.813 | 58.69050 |
| 0.9 | -259.7975 | 282.811 | 153.058 | 53.08257 |
| 0.99 | -255.8105 | 259.24 | 155.8182 | 52.66600 |
| 0.999 | -255.2795 | 260.049 | 153.2285 |  |

Table 3.2: Summary of colluders' average winnings for different positions.

The last column of Table 3.2 shows the average winnings of a pair of colluding agents playing one-million-hand matches of [2-4] Hold'em poker with different $\lambda$ values. Two colluding agents and one non-colluding agent were created by combining positional strategies. In these games, the dealer button rotates so that all of the players have equal utility gained from their position in the game. These experiments along with our previous results for positional strategies demonstrates that collusion is more beneficial when the $\lambda$ value is higher. However, the colluders must have a preference for their own utility versus their partners' when all non-colluding players fold. In such cases, setting $\lambda$ equal or very close to one would result in the colluding players acting randomly. This is shown in Table 3.2, specifically when we compare the average collusion values for $\lambda=0.9$ and $\lambda=0.999$. Therefore, $\lambda$ must be set such that colluders can differentiate between their own utility and their partners' to ensure that they behave believably when the money transitions will be between members of the colluding pair.

## Chapter 4

## Collusion Tables

In an extensive form game, every action that a player takes has the potential to affect not only on his own utility but also on the utility of all the other players in the game. As each strategy is a distribution over actions and differs from other strategies by the actions it chooses during a game, every strategy also affects the utility of players differently. Collusive strategies not only have an effect on the utility of all players in the game, but since they aim to increase the joint utility of colluders, they explicitly affect the utility of colluders in a manner different from other strategies.

Our idea for detecting collusion takes advantage of this difference. The main component of our approach is called a collusion table, which is a data structure that aims to capture the effect of players on the utility of other players. Using collusion tables, one can investigate if colluding and non-colluding strategies affect players utilities differently.

In this chapter, the semantics of a collusion table is described, followed by some examples of collusion tables. A method for creating entries in a collusion table based on a given value function is introduced. Finally, different scoring methods are presented for detecting unusual behaviours that might be signs of collusion using a collusion table.

### 4.1 Semantics

A collusion table is a data structure that is designed to capture the effect of each player's actions on all the players' utilities in a game episode. Each element $C_{g}(i, j)$
of a collusion table is called a collusion value and is meant to represent the effect of column player $j$ 's actions on the utility of row player $i$, during game episode $g$. If $C_{g}(i, j)$ is positive, it shows that player $j$ benefited player $i$ by his decision making during game episode $g$ while a negative value for $C_{g}(i, j)$ indicates that player $j$ 's actions during $g$ decreased player $i$ 's utility. Table 4.1 shows an example of a collusion table for four players. In this table, the third column represents the effect of player 3's actions on the utility of all players during the game. $C_{g}(1,3)$ is equal to 5 which means that player 3's actions favor the utility of player 1 while $C_{g}(4,3)=-3$ reveals that actions taken by player 3 had a negative impact on player 4's utility by 3 units.

Collusion values in a collusion table must satisfy the following two properties:

- The sum of the collusion values in each column of a collusion table must sum up to zero. This demonstrates that a player's actions may increase or decrease the utility of all the players but its overall effect must sum up to zero.
- The sum of the collusion values in each row of a collusion table represents the expected money that the row player can gain.

One can also define a collusion table for a dataset of game episodes $D$. Every player who participates in at least one game episode of the games in $D$ has a row and a column in this collusion table. Then, a collusion table will be generated for each game episode $g \in D$ and all $C_{g}(i, j)$ values will be computed (which will be discussed in the next section). Finally, $C_{D}(i, j)$ will be computed using $C_{g}(i, j)$ values for all game episodes $g$, in which both players $i$ and $j$ participated using the following formula:

$$
\begin{equation*}
C_{D}(i, j)=\frac{1}{|G|} \sum_{g \in G} C_{g}(i, j) \quad \text { where } \quad G=\{g \in D \mid i, j \in g\} \tag{4.1}
\end{equation*}
$$

### 4.2 Design of Collusion Tables

To build the collusion tables as described in the previous section, one must have a method for generating collusion values from game episodes. Our method focuses

| Collusion <br> Table | Player 1 | Player 2 | Player 3 | Player 4 |
| :---: | :---: | :---: | :---: | :---: |
| Player 1 | -3 | 2 | 5 | -2 |
| Player 2 | -2 | -3 | -4 | 1 |
| Player 3 | 4 | 3 | 2 | -4 |
| Player 4 | 1 | -2 | -3 | 5 |

Table 4.1: An example of a collusion table
on the change of the game value between two consecutive states for each player during a game episode. We assume that we are given a real-valued function for each agent $i$ on histories $V_{i}: H \mapsto \mathbb{R}$. We will later describe sources for this value function. Then, given a game episode $g$, the contribution of agent $j$ to the utility of agent $i$ on a terminal history $z_{g}$ is defined as:

$$
\begin{equation*}
C_{g}(i, j)=\sum_{\substack{h a \sqsubseteq z_{g} \\ p(h)=j}} V_{i}(h a)-V_{i}(h) \tag{4.2}
\end{equation*}
$$

Similarly, the impact of luck on agent $i$ 's performance is as follows:

$$
\begin{equation*}
L_{g}(i)=\sum_{\substack{h a \sqsubseteq z_{g} \\ p(h)=c}} V_{i}(h a)-V_{i)}(h) \tag{4.3}
\end{equation*}
$$

As described in Section 2.4, a player's utility can be divided into skill, luck and position. Based on the definition of $C_{g}(i, j)$ we have:

$$
\begin{equation*}
S_{g}(i)=\sum_{j \in P_{g}} C_{g}(i, j) \tag{4.4}
\end{equation*}
$$

where, instead of the player's skill, $S_{g}(i)$ now indicates the sum effect of all the players' behaviours on player $i$ 's utility. The expected utility of player $i$ in history $z_{g}$ is now as follows:

$$
\begin{equation*}
u_{g}(i)=\left[\sum_{j \in P_{g}} C_{g}(i, j)\right]+L_{g}(i)+\operatorname{Pos}_{g}\left(\phi_{g}(i)\right) \tag{4.5}
\end{equation*}
$$

Using $C_{g}(i, j)$ values, we build collusion tables and ignore the effect of luck and position on the utility of players for this study of collusion since a player's only


Figure 4.1: Example of a game episode and the value function for each state. Chance nodes are shown as triangles. In the sequence of actions / represents actions which are taken by chance.
means of colluding is their actions and they cannot use the effect of luck or position for profit. Figure 4.1 shows an example of a poker game tree. In this tree, a game episode is traced in which player 2 folded as his first action and left the game. Player

1 and 3 played until the flop cards are revealed (that is after the second chance node which is shown as a triangle). In this example the value of small blind is 5 and big blind is 10 . We denote the players actions as $f(f o l d), c$ (check or call), and $r$ (bet or raise). Figure 4.1 shows how the value of luck and the collusion values are computed during a game episode. The effect of player 1 on player 3's utility and the effect of player 3 on player 1's utility are computed as two instances.

### 4.3 Collusion Scores

Assuming that we have a collusion table which accurately reflects the influence of the agents on each other, we now focus on the problem of determining the degree of collusion exhibited by a pair of players in a collusion table. To do this, we introduce collusion scores, the functions $\theta$ from pairs of players to real numbers, where the higher a score is, the greater degree of collusion exhibited by the pair.

Collusion scores evaluate collusion tables for any suspicious patterns which might be a sign of collusion. For this step, we propose different functions $\theta$ that give scores to each pair of players in accordance with the degree of collusion illustrated by that pair.

Total Impact Score. This scoring function is designed based on the primary objective of colluding players: to increase their joint utility. Hence if two players are colluding during a game episode, their actions specifically affect their joint utility. The Total Impact score computes the total effect of the alleged colluders behaviour on their utility during a game episode $g$. It is formulated as:

$$
\begin{equation*}
\vartheta_{g}^{T I}(a, b)=\sum_{i \in\{a, b\}} \sum_{j \in\{a, b\}} C_{g}(i, j) \tag{4.6}
\end{equation*}
$$

In other words, total impact score simply sums up the four values that a pair of players gained in total, due to their own activity.

Marginal Impact Score. When colluding, players differentiate between their opponents and their partners. Therefore, we can expect a gap between a specific
player's effect on his partner and the impact he has on his opponents. Marginal Impact score is designed to examine this difference and is computed as follows:

$$
\begin{align*}
\vartheta_{g}^{M a I}(a, b)= & \left(C_{g}(b, a)-\frac{1}{|N|-2} \sum_{\substack{i \in\left\{P_{g} \\
i \notin a, b\right\}}} C_{g}(i, a)\right)+  \tag{4.7}\\
& \left(C_{g}(a, b)-\frac{1}{|N|-2} \sum_{\substack{\left.j \notin P_{g} \\
j \notin a a, b\right\}}} C_{g}(j, b)\right)
\end{align*}
$$

In the above formula, the first term computes the difference between player $a$ 's effect on player $b$ and player $a$ 's effect on other players in the game on average. The same is computed for player $b$ in the second term. Hence, the marginal impact score investigates how much a pair of players stands out based on the comparison of their behaviour toward each other and other players in the game episode.

Mutual Impact Score. This function is designed to investigate how much positive utility two players transfer to each other. The idea is that a pair of colluders should gain more mutual utility in comparison to any other pair of players in the game. This function is calculated from a game episode $g$ collusion table as follows:

$$
\begin{equation*}
\vartheta_{g}^{M u I}(a, b)=C_{g}(a, b)+C_{g}(b, a) \tag{4.8}
\end{equation*}
$$

Minimum Impact Score. The fourth scoring function is based on the idea that both players must be colluding for collusion to really take place. In this measure, the possibility of collusion is investigated based on a player's individual effect on their partnership rather than the total utility that they gain together (i.e. in total impact score). This function is formulated as:

$$
\begin{equation*}
\vartheta_{g}^{M i I}(a, b)=\min _{i \in\{a, b\}}\left\{\sum_{j \in\{a, b\}} C_{g}(j, i)\right\} \tag{4.9}
\end{equation*}
$$

For each pair of players, this measure reflects the minimum utility that each partner contributed to their joint utility. In this way, the pair of players with the maximum minimum contribution from both of the partners would be ranked first. Therefore, accidental collusion in which the joint utility increases by just one of the partners would not be ranked highly.

|  | Total <br> Impact | Marginal <br> Impact | Mutual <br> Impact | Min <br> Impact | Differential <br> Total Impact |
| :---: | :---: | :---: | :---: | :---: | :---: |
| Player 1-Player 2 | -6 | -3 | 0 | -5 | -14 |
| Player 1-Player 3 | 8 | 13 | 9 | 1 | 7 |
| Player 1-Player 4 | 1 | -0.5 | -1 | -2 | -7 |
| Player 2-Player 3 | -2 | -2 | -1 | -2 | -10 |
| Player 2-Player 4 | 1 | -0.5 | -1 | -5 | -7 |
| Player 3-Player 4 | 0 | -7 | -7 | -1 | -8 |

Table 4.2: Different collusion scores for a sample collusion table shown in Table 4.1.

Differential Score. This scoring method tries to determine the difference between the score of the most suspicious pair and the rest of the pairs of players in the game. This method can be defined using any other measure of collusion as its base score. Here, we use the Total Impact score as the base scoring method. Differential Total Impact score is computed as follows:

$$
\begin{equation*}
\vartheta_{g}^{D I}(a, b)=\vartheta_{g}^{T I}(a, b)-\max \left\{\max _{\substack{d \in P_{g} \\ d \notin\{a, b\}}}\left\{\vartheta_{g}^{T I}(a, d)\right\}, \max _{\substack{d \in P_{g} \\ d \notin\{a, b\}}}\left\{\vartheta_{g}^{T I}(b, d)\right\}\right\} \tag{4.10}
\end{equation*}
$$

Using this method, the top value among the base scores stands out and all the rest of the pairs of players' values will be negated. Therefore, this method can be used when there is a need to detect the most suspicious pair of players rather than to order the pairs of agents according to their likeliness of collusion.

Table 4.2 demonstrates the different scoring methods for the collusion table example shown in Table 4.1. Based on all of the measures, player 1 and player 3 have the maximum collusion scores which shows that their behaviour is the most suspicious for collusion. The data structure and measurement methods that were described in this chapter provide the basis for our collusion detection system which is experimentally verified in Chapter 7.

## Chapter 5

## Creating Value Functions for Sequential Games

To create collusion tables, a method of determining the value for each agent in every history in the game is needed. This value function will be used to determine the effect of each player's actions on all the players' utilities.

We have already seen the similarity between the underlying methods of the advantage sum estimator described in Chapter 2 and the generation of collusion values in a collusion table. These methods can utilize any given value function. Some advantage sum estimators use a hand crafted domain specific value function (like DIVAT [2]) while others use a learned value function using features of the domain (like MIVAT [23]).

In this chapter, three value functions are introduced: Always call, Purified CFR, and PIVAT. The first two methods take advantage of a base strategy to implicitly define the value function, while the last method learns the value function based on the features of the domain. Note that each of these value functions are learned for each position (player) in the game. Then, a set of value functions is used to evaluate a specific agent in different positions.

### 5.1 Value Functions Based on a Strategy

Any strategy can be used as the basis of defining a value function for a game. The value function then is simply the expected utility gained from any game state for each player if all the players were to use the base strategy to make their decisions
during the remainder of the game episode. If a good non-colluding strategy is used as a value function and it reveals that the utility gained by a group of players is unexpectedly high, those players are more likely to be colluding. In this section, two strategies are described that are used in our experiments as a value function. The first one is a very simple strategy while the second one is one of the state of the art strategies that have been shown to have good performance in large extensive form games.

### 5.1.1 Always Call

The first strategy that we utilized to define a value function is called always call. This strategy selects the "call" action in all game states and for all players without considering the private or public cards. In order to compute the value of the game in a history for a player $i$, it assumes that all the participants would call for the rest of the game. This strategy is naive but has been shown in other domains to yield significant reduction in variance [18].

### 5.1.2 Purified CFR

The second value function used in this work is defined by a strategy that is generated using Counterfactual Regret Minimization method (CFR), a no-regret learning algorithm $[16,28]$. In order to increase the speed of evaluation in each game state, we de-randomized the base CFR strategy by deterministically choosing the action assigned the highest probability by the CFR base strategy in each information set, a technique which is introduced by Ganzfried et al. [5]. This value function is named purified CFR.

### 5.2 Implementation of Value Functions

The Always call value function was implemented enumerating all the possible outcomes at a chance node. As an example, a small experiment is performed consisting of 3 positional agents which play using identical non-colluding CFR strategies for a set of one million hands. Table 5.1 shows the collusion table for these games, as

|  | Agent 1 | Agent 2 | Agent 3 |
| :---: | :---: | :---: | :---: |
| Agent 1 | 332.389024 | -214.061441 | -213.230632 |
| Agent 2 | -243.602725 | 420.087003 | -358.769791 |
| Agent 3 | -88.786299 | -206.025562 | 572.000423 |

Table 5.1: Collusion table for 3 non-colluding positional agents evaluated by always call policy ( $\mathrm{mbb} / \mathrm{g}$ ).

|  | Agent 1 | Agent 2 | Agent 3 |
| :---: | :---: | :---: | :---: |
| Agent 1 | -1.08161 | -0.274575 | 0.443451 |
| Agent 2 | 0.785583 | 0.804984 | 1.06696 |
| Agent 3 | 0.296089 | -0.530411 | -1.51039 |

Table 5.2: Collusion table for 3 non-colluding positional agents evaluated by determinized CFR policy ( $\mathrm{mbb} / \mathrm{g}$ ).
evaluated using the always call defined value function.
We also implemented a simplified version of the Purified CFR value function to speed up the evaluation process. In this implementation, to determine the value of a chance node, only a single chance outcome is sampled, instead of enumerating all outcomes as with the always call value function. We also evaluated another set of one million game episodes using purified CFR. In these games, 3 positional agents played using identical non-colluding CFR strategies. Table 5.2 shows the resulting collusion table for this set of games.

The values in Table 5.1 are rather large, demonstrating that the strategy that players actually used in the game episodes was quite different from the policy used to define the value function. In such a case, the policy would compute an expected value of player $i$ for a history $h$ based on the strategy used as the value function. The player then would take action $a$ and reach history $h a$ while the value function assumed that it would choose another action $a^{\prime}$ and reaches $h a^{\prime}$. Therefore, the collusion value computed from the values of the game in histories $h$ and ha may become large. On the other hand, Table 5.2 shows much lower values. This is due to the greater degree of similarity between the value function strategy and the strategies used by the players in the game episodes.

### 5.3 PIVAT

The two policies discussed earlier use a base strategy. The quality of evaluation depends on how close the base strategy and the players' strategies are. Therefore, a better method may be to use a dataset of games previously played by agents and learn a value function for that population. Using such a value function improves collusion detection technique since it is tailored to the target population and can also reduce the variance as shown by White and Bowling [23].

The third value function examined in this thesis is called the PIVAT value function which learns an estimator of collusion values. This method is based on similar ideas from the DIVAT and MIVAT assessment tools. In this approach, a value function is designed for each player based on a set of features for the game.

### 5.3.1 Finding the Value Function

The main goal for the PIVAT value function is to minimize the variance of collusion values so as to better estimate the effect of players' actions on the utilities of all players. Assume we are given a value function $V_{i}: H \mapsto \mathbb{R}$ for each player $i$. We first reformulate the impact of player $j$ 's actions on player $i$ 's utility during game episode $g$ as:

$$
\begin{equation*}
C_{g}(i, j)=1\left[P_{-1}\left(z_{g}\right)=j\right]\left(u_{i}\left(z_{g}\right)-V_{i}\left(z_{g}\right)\right)+\sum_{\substack{h a \sqsubseteq z_{g} \\ P(h)=j}} V_{i}(h a)-V_{i}(h) \tag{5.1}
\end{equation*}
$$

in which $P_{-1}\left(z_{g}\right)$ indicates the player who made the last move before the game reached terminal history $z$. In this formula, the indicator function $\left(1\left[P_{-1}\left(z_{g}\right)=j\right]\right)$ is equal to 1 when the last player was player $j$, otherwise it equals 0 .

Note that the variance of collusion values will depend on the strategies which players use in the game episodes. Our goal is to learn a better value function using the past interactions of players. Therefore, to estimate the variance, some samples of game episodes $G=\left\{g_{1}, g_{2}, \ldots, g_{T}\right\}$ are utilized. Our objective is then to solve the following optimization problem given a dataset of game episodes $G$ :

$$
\begin{equation*}
\min _{V_{i} \in \mathbb{V}_{i}} \sum_{j \in N} \operatorname{Var}\left(C_{G}(i, j)\right) \tag{5.2}
\end{equation*}
$$

where $\mathbb{V}_{i}$ is the class of all value functions we are considering. The variance of a collusion value in a game episode $g$ can be written as:

$$
\begin{align*}
\operatorname{Var}\left(C_{G}(i, j)\right) & =\mathbb{E}_{g \sim G}\left[\left(C_{G}(i, j)-\bar{C}(i, j)\right)^{2}\right] \\
& \approx \frac{1}{T} \sum_{t=1}^{T}\left[\left(C_{g_{t}}(i, j)-\frac{1}{T} \sum_{t^{\prime}=1}^{T} C_{g_{t^{\prime}}}(i, j)\right)^{2}\right] \tag{5.3}
\end{align*}
$$

To ensure that the estimator using this value function satisfies the properties of a collusion table, the zero-luck constraint must be satisfied. Thus means that $E\left[L_{V_{j}} \mid \sigma\right]=0$. This constraint is equivalent to requiring that for each player, the value of a state right before chance is to act must be equal to the expected value of the states right after the chance node. If this constraint is satisfied, the expected money gained by a player $i$ will be equal to the sum effect of all the players' behaviour on player $i$ 's utility. As explained in Chapter 2, in order to satisfy this constraint, we define the value of a chance node as follows:

$$
\begin{equation*}
V_{i}(h \text { s.t. } P(h)=c)=\sum_{a \in A(h)} \sigma_{c}(a \mid h) V_{i}(h a) \tag{5.4}
\end{equation*}
$$

This guarantees that the value function will be unbiased.

### 5.3.2 Linear Value Functions

We focus on the class of linear value functions, since we want the optimization to be tractable. Define a feature mapping $\lambda: H \mapsto \mathbb{R}^{n}$ which maps histories to a vector of $n$ real-valued features. We then consider value functions which are a linear combination of these features:

$$
V_{i}(h)=\theta_{i}^{T} \lambda(h)
$$

for some $\theta_{i} \in \mathbb{R}^{n}$. We can now rewrite the definition of collusion values using the new value function.

$$
\begin{equation*}
C_{g}(i, j)=\mathbf{1}\left[P_{-1}\left(z_{g}\right)=j\right]\left(u_{i}\left(z_{g}\right)-\theta_{i}^{T} \lambda(z)\right)+\theta_{i}^{T}\left(\sum_{\substack{h a \sqsubset z \\ P(h)=j}} \lambda(h a)-\lambda(h)\right) \tag{5.5}
\end{equation*}
$$

additionally, the zero-luck constraint can be rewritten as:

$$
\lambda(h)=\sum_{a} \sigma_{c}(a \mid h) \lambda(h a)
$$

### 5.3.3 Linear Optimization Objective

In this section, the full optimization objective function is written out and then simplified. Given a dataset of game episodes $G=\left\{g_{1}, g_{2}, \ldots, g_{t}\right\}$ and for each player $i \in N$, our goal is:

$$
\begin{align*}
& \underset{\theta_{i} \in \mathbb{R}^{n}}{\operatorname{Minimize}:} \sum_{j \in N}\left\{\frac { 1 } { | G | } \sum _ { g \in G } \left[\theta_{i}^{T}\left(\sum_{\substack{h a \sqsubseteq z_{g} \\
P(h)=j}} \lambda(h a)-\lambda(h)\right)\right.\right. \\
& \quad+\mathbf{1}\left[P_{-1}\left(z_{g}\right)=j\right]\left(u_{i}\left(z_{g}\right)-\theta_{i}^{T} \lambda\left(z_{g}\right)\right) \\
& \quad-\frac{1}{|G|} \sum_{g^{\prime} \in G}\left[\theta_{i}^{T}\left(\sum_{\substack{h a \sqsubseteq z_{g^{\prime}} \\
P(h)=j}} \lambda(h a)-\lambda(h)\right)\right. \\
& \left.\left.\left.\quad+\mathbf{1}\left[P_{-1}\left(z_{g^{\prime}}\right)=j\right]\left(u_{i}\left(z_{g^{\prime}}\right)-\theta_{i}^{T} \lambda\left(z_{g^{\prime}}\right)\right)\right]\right]^{2}\right\} \tag{5.6}
\end{align*}
$$

This expression can be simplified as follows

$$
\begin{array}{r}
\sum_{j \in N}\left\{\frac { 1 } { | G | } \sum _ { g \in G } \left[\theta_{i}^{T}\left(-\mathbf{1}\left[P_{-1}\left(z_{g}\right)=j\right] \lambda\left(z_{g}\right)+\sum_{\substack{h a \sqsubseteq z_{g} \\
P(h)=j}} \lambda(h a)-\lambda(h)\right)\right.\right. \\
+\mathbf{1}\left[P_{-1}\left(z_{g}\right)=j\right]\left(u_{i}\left(z_{g}\right)\right) \\
-\frac{1}{|G|} \sum_{g^{\prime} \in G}\left[\theta_{i}^{T}\left(-\mathbf{1}\left[P_{-1}\left(z_{g^{\prime}}\right)=j\right] \lambda\left(z_{g^{\prime}}\right)+\sum_{\substack{h a \sqsubseteq z_{g^{\prime}} \\
P(h)=j}} \lambda(h a)-\lambda(h)\right)\right. \\
 \tag{5.7}\\
\left.\left.\left.+\mathbf{1}\left[P_{-1}\left(z_{g^{\prime}}\right)=j\right]\left(u_{i}\left(z_{g^{\prime}}\right)\right)\right]\right]^{2}\right\}
\end{array}
$$

By defining the following shorthand notation,

$$
A_{g}(i, j)=-\mathbf{1}\left[P_{-1}\left(z_{g}\right)=j\right] \lambda\left(z_{g}\right)+\sum_{\substack{h a \sqsubseteq z_{g} \\ P(h)=j}} \lambda(h a)-\lambda(h)
$$

$$
\begin{array}{r}
\begin{array}{r}
\bar{A}(i, j)=\frac{1}{|G|} \sum_{g^{\prime} \in G}\left[\left(-\mathbf{1}\left[P_{-1}\left(z_{g^{\prime}}\right)=j\right] \lambda\left(z_{g^{\prime}}\right)\right)+\sum_{\substack{h a \sqsubseteq z_{g^{\prime}} \\
P(h)=j}} \lambda(h a)-\lambda(h)\right] \\
=\frac{1}{|G|} \sum_{g^{\prime} \in G} A_{g^{\prime}}(i, j) \\
B_{g}(i, j)=\mathbf{1}\left[P_{-1}\left(z_{g}\right)=j\right] u_{i}\left(z_{g}\right)-\frac{1}{|G|} \sum_{g^{\prime} \in G} \mathbf{1}\left[P_{-1}\left(z_{g^{\prime}}\right)=j\right] u_{i}\left(z_{g^{\prime}}\right)
\end{array} .
\end{array}
$$

we get the following optimization:

$$
\underset{\theta_{i} \in \mathbb{R}^{n}}{\operatorname{Minimize}} \sum_{j \in N} \sum_{g \in G}\left[\theta_{i}^{T}\left(A_{g}(i, j)-\bar{A}(i, j)\right)+B_{g}(i, j)\right]^{2}
$$

### 5.3.4 Linear Optimization Solution

As our objective is convex in $\theta_{i}$, we can solve this optimization by setting the objective's partial derivative to 0 .

$$
\nabla_{\theta_{i}} J\left(\theta_{i}\right)=\nabla_{\theta_{i}}\left[\sum_{j \in N} \sum_{g \in G}\left[\theta_{i}^{T}\left(A_{g}(i, j)-\bar{A}(i, j)\right)+B_{g}(i, j)\right]^{2}\right]
$$

$$
\begin{gathered}
\nabla_{\theta_{i}} J\left(\theta_{i}\right)=2 \sum_{j \in N} \sum_{g \in G}\left[A_{g}(i, j)-\bar{A}(i, j)\right]\left[\theta_{i}^{T}\left(A_{g}(i, j)-\bar{A}(i, j)\right)+B_{g}(i, j)\right] \\
\nabla_{\theta_{i}} J\left(\theta_{i}\right)=\left[\sum_{j \in N} \sum_{g \in G}\left(A_{g}(i, j)-\bar{A}(i, j)\right)\left(A_{g}(i, j)-\bar{A}(i, j)\right)^{T}\right] \theta_{i} \\
+\sum_{j \in N} \sum_{g \in G}\left(A_{g}(i, j)-\bar{A}(i, j)\right) B_{g}(i, j)
\end{gathered}
$$

if we set $\nabla_{\theta_{i}} J\left(\theta_{i}\right)=0$ and solve for $\theta_{i}^{*}$

$$
\begin{aligned}
& \theta_{i}^{*}=\left[\sum_{j \in N} \sum_{g \in G}\left(A_{g}(i, j)-\bar{A}(i, j)\right)\left(A_{g}(i, j)-\bar{A}(i, j)\right)^{T}\right]^{-1} \\
& \times {\left[\sum_{j \in N} \sum_{g \in G}\left(\bar{A}(i, j)-A_{g}(i, j)\right) B_{g}(i, j)\right] }
\end{aligned}
$$

The linear value function defined by these weights is the PIVAT value function. The experimental results in this thesis will focus on the other value function methods though. We leave experimental evaluation of PIVAT for future work.

## Chapter 6

## Experimental Design

### 6.1 Introduction

In the previous chapters, we described our collusion detection method. To evaluate the effectiveness of this method, a dataset of games played by different types of players is needed. The most desirable dataset for evaluating collusion detection methods would be one composed of real-world, labelled human data. However, due to the unavailability of such a dataset, we designed and constructed a synthetic dataset of poker games with different types of agents. In this chapter, different strategies and the abstractions used to create weak and strong strategies are explained. Finally, the different types of agents that used our designed strategies to create the dataset of games are introduced.

### 6.2 Creating Different Strategies

The next step in building an experimental platform is to design and create agents. Each agent is constructed by combining a set of positional strategies (one positional strategy for each position in the game). In these experiments, all agents are designed to play 3-player poker games since it is the smallest multi-player poker game in which collusion is possible. Every agent employs a set of strategies to play poker. Positional strategies were built using the Chance Sampling CFR algorithm (a variant of CFR which samples one set of chance outcomes per iteration) [12,28].

Due to the large size of the poker game tree, strategies are built in an abstracted version of the game. An abstracted game shares the same strategic properties with
the full game while having fewer game states and information sets [9]. During the running time of the CFR algorithm, the strategies for the players are updated during a set of self-play games. Therefore, in each run of CFR, 3 positional strategies will be generated (that is: one for each position in the game). First, an abstracted game tree is built and initialized for each player in the game. Then, in each iteration and for each information set, the action probabilities for that player are calculated using the accumulated counterfactual regret of playing each possible action. Finally, due to the large number of choices in a chance node, a sampling method will be used instead of enumerating all the possible outcomes. In general, the higher the number of iterations are, the closer to optimal the developed positional strategy will be.

Each created strategy is position dependent and also created in the context of a specific situation, depending on whether the other players are its partners, colluding opponents or normal players. An agent is constructed from three positional strategies. We developed three different kinds of strategies to construct our agents from:

- Colluder. This kind of strategy is built using the modified utility function that is described in Chapter 3. For each combination of two out of three positions, a pair of colluding strategies is developed.
- Defender. As mentioned before, during iterations of the CFR algorithm a strategy is developed for each position in the game. Therefore, when a pair of positional colluding strategies are being generated, the algorithm also produces a third strategy for the third position. This strategy plays to minimize its regret while playing against colluding opponents. We call this strategy the defender.
- Normal. When the CFR algorithm is run without modifying the players' utility functions, three positional strategies are developed which are playing unmodified [2-4] Holdem. These strategies minimize their regret assuming that their opponents are also normal players.


### 6.2.1 Abstractions

Computing a strategy in a large extensive form game like poker usually requires an enormous amount of memory. Although [2-4] Hold'em is a smaller variation of Texas Hold'em, it still needs a huge amount of computer memory to build and store a strategy for it since it has $\sim 10^{17}$ game states and $1.677 \times 10^{12}$ information sets. Therefore, abstraction is utilized to reduce the size of the game tree. An abstracted version of the game is created combining similar information sets into one bucket. A strategy will be created then in this abstracted game.

Beside managing the memory size issue, we also use abstractions as a tool to create strategies with different strengths. Based on the method of abstraction and the number of buckets used, the resulting strategies will perform with different strength in the real game. In these experiments, we used abstractions which vary both in size and in the method of abstraction to develop weak and strong strategies. The stronger agents were created using a bigger abstraction while smaller abstractions were utilized to develop strategies for weak agents. These abstractions are introduced below:

- 3700-K-Means Abstraction. In this abstraction, complete knowledge about the pre-flop is used. That means there is one information set in the abstracted game for every information set in the real game. On the flop round, information sets are grouped into 3700 different buckets. K-means clustering is used to create buckets based on hand strength. This abstraction uses imperfect recall, meaning that the exact pre-flop cards may be forgotten on the flop. We utilized this abstraction to construct strong agents.
- Ns-Bucket Abstractions. This abstraction divides the information sets in each round into N fixed buckets based on the hand strength information. These abstractions use perfect recall in contrast with the previous method of abstraction. Two sets of weak agents were developed using 2-bucket and 5-bucket abstractions. The first set of agents can distinguish between 2 buckets on the pre-flop round and 4 buckets on the flop. Similarly, the second set can distinguish between 5 buckets on the pre-flop round and 25 buckets on
the flop which is extremely small in comparison with the actual number of information sets. As a result, strategies generated in this abstracted game may play weakly in the real game since they must use same action probabilities for all hands in a bucket, despite the individual hands having considerable differences.


### 6.3 Corpus of Games

Using the strategies and abstractions that are described in the previous section, we designed agents to form a population of diverse participants. In a real ecosystem of poker games such as an online poker game, users participate with various abilities. We created seven different types of agents by combining the positional strategies described earlier.

- Colluder A. (CA) This agent colludes with his partner colluder B when both of them are in a game. When this occurs, Colluder A applies an appropriate Colluder strategy based on his seat and his partner's seat in each game episode. If Colluder B is not one of the players participating in the game episode, Colluder A utilizes a Normal strategy.
- Colluder B. (CB) This player colludes with Colluder A when both of them are participating in a game and plays symmetrically.
- Non-Colluder. (NC) This agent applies a Normal strategy in all seats during a game with no dependence on who his opponents are.
- Defender. (DF) In a real game situation, there may be agents who are suspicious that their opponents are colluding and hence try to play their best response to collusion. The defender agent is designed to represent the group of smart agents who can detect and respond correctly if their opponents are colluding. This agent employs the Defender strategy in such situations. Otherwise, it uses a Normal strategy.
- Paranoid. (PR) While Defender agents can always detect occurrence of collusion correctly, there are other agents who are suspicious that their opponents are colluding. They always play a defensive strategy while they may be right or wrong. The Paranoid agent is designed to represent this group and it always employs a Defender strategy regardless of its opponents.
- Accidental collude-right. (CR) It is possible that agents are behaving as if they are colluding due to their lack of skill and this can benefit other players in a real game situation. The accidental colluding agents are designed to represent such implicit acts of collusion. This agent always utilizes a colluder strategy that benefits the agent on his right, regardless of that agent's strategy.
- Accidental collude-left. (CL) Similar to the Accidental collude-right agent, this agent always employs a colluder strategy designed to collude with the agent on his left.

Three sets of agents were created: one strong set and two weak sets each consisting of seven agents using $\lambda=0.9$. Strong agents name are in the form of S.x and weak agents are in the form of WN.x in which N is the number of buckets used in their abstraction and $x$ indicates the type of the agent. Every combination of three agents out of fourteen (one strong set and one weak set) played a one-million-hand match (364 sets in total).

## Chapter 7

## Results and Discussion

In this chapter we present the results of our collusion detection system, run with two of the proposed collusion value functions, for each of our test populations of agents. The results of experiments using the always-call value function are presented first. We do not include the complete experiments for this value function since it does not perform well. Then, the experiments using the purified CFR value function are described. Collusion tables were created for each game episode and combined into one collusion table using formula 4.1 for each match. Then collusion values for each pair of agents obtained in all of the matches in which both had participated were averaged and merged into one collusion value. Finally, we evaluated the final collusion table using all of the collusion scoring methods introduced in Chapter 5. The resulting set of rankings for pairs of agents are presented in this chapter.

### 7.1 Using the Always-Call Value function

Two one-million-hand matches were played by strong agents to evaluate the performance of our collusion detection method using the always-call value function described in Chapter 5. The first set was played by two colluding agents and one non-colluding agent. The second set was played by two colluding agents and one defending agent. The collusion tables created from these games using the Alwayscall value function are presented in Tables 7.1 and 7.2 and the resulting ranking from all scoring methods are shown in Tables 7.3 and 7.4.

None of the the scoring methods except the Minimum Impact score could detect

| Agent | S.CA | S.CB | S.NC |
| :---: | :---: | :---: | :---: |
| S.CA | 480.653 | -208.739 | -171.957 |
| S.CB | -280.891 | 409.611 | -175.271 |
| S.NC | -199.761 | -200.872 | 347.228 |

Table 7.1: Match1: Collusion table for 3 agents in mbb/g using the always-call value function.

| Agent | S.CA | S.CB | S.DF |
| :---: | :---: | :---: | :---: |
| S.CA | 461.828 | -201.985 | -164.893 |
| S.CB | -276.89 | 388.476 | -183.15 |
| S.DF | -184.938 | -186.49 | 348.044 |

Table 7.2: Match2: Collusion table for 3 agents in mbb/g using always-call value function.

| Agent $i$ | Agent $j$ | Total <br> Impact | Mutual <br> Impact | Marginal <br> Impact | Minimum <br> Impact | Differential Total <br> Impact |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| S.CA | S.NC | 456.163 | -371.718 | 84.444 | 175.271 | 55.529 |
| S.CA | S.CB | 400.634 | -489.63 | -88.997 | 199.762 | -55.529 |
| S.CB | S.NC | 380.696 | -376.143 | 4.553 | 171.957 | -75.467 |

Table 7.3: Result of different scoring methods for Match1.

| Agent $i$ | Agent $j$ | Total <br> Impact | Mutual <br> Impact | Marginal <br> Impact | Minimum <br> Impact | Differential Total <br> Impact |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: |
| S.CA | S.DF | 460.041 | -349.831 | 110.209 | 183.151 | 88.612 |
| S.CA | S.CB | 371.429 | -478.875 | -107.447 | 184.938 | -88.612 |
| S.CB | S.DF | 366.88 | -369.64 | -2.762 | 164.894 | -93.161 |

Table 7.4: Result of different scoring methods for Match2.
the colluding agents in either matches, even among only three pairs of agents. Also, the Minimum Impact score of the colluders is only a little bit higher than the next highest score which shows that it does not distinguish colluding agents reliably. These results show that, even though it has been shown to have good performance in agent evaluation, the always-call value function does not have good performance when used for collusion detection.

### 7.2 Strong vs. 5s Weak Agents Using Purified CFR Value Function

Since the always-call value function did not work, this led us to develop the Purified CFR value function described in Chapter 5, which we will show, can be used to detect collusion. We now describe the results of matches between strong agents and weak agents that were built based on 5 s abstractions. The value function used in these experiments is Purified CFR. In these experiments all of the scoring methods ranked both the strong and weak colluders at or near the top of the rankings. The complete ranking of agents in this experiment is presented in the first section of the Appendix.

### 7.2.1 Money Winnings

Table 7.5 shows the top twelve pairs of agents ranked based on the money they gained in all the matches in which both participated. This table shows that the strong colluding agents (S.CA and S.CB) gained the most money when playing together, which is a very strong sign that they may have been colluding. However, it does not clearly identify the weak colluding agents (W5.CA and W5.CB) as they are ranked in the twelfth position. Also, two strong agents in a field of generally weak agents may appear to collude. Therefore, using the total money that pair of agents gained can be a sign of collusive behaviour but it has both false positive and false negative results and is not enough to always identify collusion. Figure 7.1 shows the distribution of agent pairs based on the total money they gained across all matches. This histogram reveals that collusion was beneficial for both pairs
of colluders as even the weak colluding agents gained a considerable amount of money.

| Rank | Agent $i$ | Agent $j$ | Money gained | Rank | Agent $i$ | Agent $j$ | Money gained |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | S.CA | S.CB | 86.359 | 7 | S.PR | S.CB | 30.483 |
| 2 | S.DF | S.CA | 39.023 | 8 | S.PR | S.CA | 30.063 |
| 3 | S.DF | S.NC | 38.920 | 9 | S.PR | S.NC | 29.968 |
| 4 | S.CA | S.NC | 37.228 | 10 | S.DF | S.PR | 29.544 |
| 5 | S.DF | S.CB | 36.648 | 11 | W5.CA | W5.CB | 24.036 |
| 6 | S.NC | S.CB | 36.560 | 12 | S.CL | S.DF | 17.747 |

Table 7.5: Experiment 1: Top twelve ranking of pairs of agents based on money gained in mbb/g.


Figure 7.1: Experiment 1: Money gained by pairs of agents histogram.

### 7.2.2 Total Impact Score

The ranking of agent pairs based on their Total Impact score is presented in Table 7.6. In this table, the strong and weak colluders are ranked as the first and second most likely pairs of agents to be colluding. The strong colluders score is very high when compared with the rest of the agent pair scores. This difference can be seen in the distribution of scores which is shown in Figure 7.2. This shows that not only are both colluders at the top of the ranking, but the strong colluders are clear outliers.

| Agent $i$ | Agent $j$ | TI |
| :---: | :---: | :---: |
| S.CA | S.CB | 1.00555 |
| W5.CA | W5.CB | 0.385842 |
| S.CR | W5.CR | 0.375167 |
| S.CL | S.CR | 0.349336 |
| W5.CL | S.CR | 0.091698 |

Table 7.6: Experiment 1: Top five ranking of pairs of agents based on Total Impact score.


Figure 7.2: Experiment 1: Total Impact Score histogram.

| Agent $i$ | Agent $j$ | MaI |
| :---: | :---: | :---: |
| W5.CR | S.CR | 1.19568 |
| S.CA | S.CB | 1.09196 |
| W5.CA | W5.CB | 0.454279 |
| S.CL | S.CR | 0.219464 |
| S.CL | W5.CR | 0.0536221 |

Table 7.7: Experiment 1: Top five ranking of pairs of agents based on Marginal Impact Score.


Figure 7.3: Experiment 1: Marginal Impact Score histogram.

### 7.2.3 Marginal Impact Score

Table 7.7 shows the top five ranking agent pairs based on the Marginal Impact score. Based on this measure, strong and weak colluding agents are ranked second and third respectively which shows that this method also can detect colluders. Strong and weak accidental colluders-right compose the pair of agents with the highest score in this table. The top four rankings seem to be outliers in Figure 7.3. Looking at Table 7.7 one can see that the agent pairs in the top five spots are either colluding agents or accidental colluders. These results show that although they are not intentionally colluding, their behaviour is suspicious. Our method with the Marginal Impact score detects this and further investigation by human experts would be needed to ensure they are not colluding.

### 7.2.4 Mutual Impact Score

Table 7.8 shows the top eighteen ranking pairs of agents based on the Marginal Impact score. As shown in this table and Figure 7.4, this measure does not appear to adequately detect collusion. This experiment demonstrates that although the colluding agents gain more money in total, our system does not detect that they are transferring utility to each other significantly more than other pairs of agents.

| Rank | Agent $i$ | Agent $j$ | MuI | Rank | Agent $i$ | Agent $j$ | MuI |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | S.CR | W5.CR | 0.8205 |  | 10 | S.CB | W5.CR |
|  | 0.0828874 |  |  |  |  |  |  |
| 2 | W5.CR | W5.NC | 0.177952 | 11 | W5.CR | W5.PR | 0.0819924 |
| 3 | W5.CR | W5.DF | 0.17314 | 12 | S.DF | W5.PR | 0.0799518 |
| 4 | W5.CA | W5.CR | 0.158265 | 13 | W5.NC | W5.PR | 0.0775214 |
| 5 | W5.CB | W5.CR | 0.148728 | 14 | S.CA | W5.CR | 0.0747294 |
| 6 | W5.CR | S.DF | 0.145593 | 15 | S.NC | W5.PR | 0.0720325 |
| 7 | W5.CR | S.NC | 0.14041 | 16 | W5.DF | W5.PR | 0.0711563 |
| 8 | S.PR | W5.PR | 0.0898325 | 17 | W5.CA | W5.CB | 0.0684351 |
| 9 | S.CA | S.CB | 0.0864108 | 18 | W5.DF | W5.NC | 0.0651412 |

Table 7.8: Experiment 1: Top eighteen ranking of pairs of agents based on Mutual Impact Score.


Figure 7.4: Experiment 1: Mutual Impact Score histogram.

### 7.2.5 Minimum Impact Score

Table 7.9 shows that colluders are close to top of the ranking based on Minimum Impact score. Again, using this measure, both pairs of colluding agents can be distinguished. One nice feature of this scoring function, as shown in this table and also Figure 7.5, is that only the top four pairs have positive scores among all 91 pairs and both colluding pairs are among those. So, this measure appears to attach meaning to the sign of a pair's collusion score.

| Agent $i$ | Agent $j$ | MiI |
| :---: | :---: | :---: |
| S.CA | S.CB | 0.414292 |
| S.CL | S.CR | 0.0964158 |
| W5.CA | W5.CB | 0.0406547 |
| S.DF | S.NC | 0.00535241 |
| S.NC | S.CB | -0.00691819 |

Table 7.9: Experiment 1: Top five ranking of pairs of agents based on Minimum Impact Score.


Figure 7.5: Experiment 1: Minimum Impact Score histogram.

| Agent $i$ | Agent $j$ | DI |
| :---: | :---: | :---: |
| S.CA | S.CB | 0.59348 |
| W5.CR | S.CR | 0.281338 |
| W5.CA | W5.CB | 0.262688 |
| S.CL | S.CR | 0.0502188 |
| S.DF | S.PR | 0.00218301 |

Table 7.10: Experiment 1: Top five ranking of pairs of agents based on Differential Total Impact Score.


Figure 7.6: Experiment 1: Differential Total Impact Score histogram.

### 7.2.6 Differential Total Impact Score

The top five ranking pairs of agents based on the Differential Total Impact score are presented in Table 7.10. The strong colluders are detected as the most suspicious pair based on this measure. The complete ranking of pairs, presented in Table A.5, shows that only the first five values are positive. As explained before, this measure tries to magnify the difference between the most suspicious pairs and the rest of the pairs of agents. Figure 7.6 demonstrates that it was successful at this goal. The two colluding pairs are ranked in the first and third position. Also the accidental colluding agents appeared in the second and fourth position which again shows that their behaviour is suspicious and requires further investigation.

### 7.3 Strong vs. 2s Weak Agents Using Purified CFR Value Function

In this section, the result of a set of matches between strong agents and weak agents that were built using a 2 s abstractions (which is described in Chapter 6) are presented. Each of the weak agents that were used in these experiments only had one bit of information (good/bad) about each chance event (i.e. good/bad private cards, good/bad flop cards given these private cards). The value function used in these experiments is Purified CFR. The complete ranking of agents in this experiment are presented in the second section of the Appendix.

Figure 7.7 demonstrates that the strong colluders gained the most money and appeared as an outlier among the other pairs of agents. However, the weak colluding agents did not gain a considerable amount of money. Table A. 7 shows that not only did the weak colluders not gain money but also they lost money by $-19.0389 \mathrm{mbb} / \mathrm{g}$ and were ranked in sixtieth place. This shows that their colluding strategies were not particularly beneficial.


Figure 7.7: Experiment 2: Money gained by pairs of agents histogram.

The histograms in Figures 7.8-7.12 show that the strong colluders were detected by all scoring methods in the first or second place except for Mutual Impact score. This experiment also supports our previous idea that using, Mutual Impact score


Figure 7.8: Experiment 2: Total Impact Score histogram.
does not well identify colluding agents. However, other scoring methods detected the strong colluders as an obvious outlier when compared with the rest of the pairs' scores.

None of the scoring methods could detect the weak colluders. We hypothesize that this is due to the weak abstractions that the weak colluding agents utilized. Their collusion did not even yield them a profit, so in some sense, they were attempting to collude, but they were not successful. These agents resemble beginner players who try to collude in poker but they just harm each other by trying to collude. Figures 7.8-7.12 show that our scoring methods cannot detect this type of colluders, or rather, those who fail to successfully collude.


Figure 7.9: Experiment 2: Marginal Impact Score histogram.


Figure 7.10: Experiment 2: Mutual Impact Score histogram.


Figure 7.11: Experiment 2: Minimum Impact Score histogram.


Figure 7.12: Experiment 2: Differential Total Impact Score histogram.

### 7.4 Discussion

The experiments presented in this chapter show that our collusion detection system using the purified CFR value function can successfully detect collusion if collusion is beneficial for colluding agents. Table 7.11 demonstrates a summary of the performance of the scoring methods on different datasets. It shows that all the scoring methods except the Mutual Impact score can detect strong colluding agents when Purified CFR is used as the value function. However, the Always-call value function does not show promising performance.

All of the scoring methods introduced in this thesis can be used in other multiagent domains. Therefore, we recommend using a combination of these scoring methods in order to have a robust collusion detection method.

|  | Purified CFR value function |  | Always-call |
| :---: | :---: | :---: | :---: |
| Scoring method | Experiment 1 | Experiment 2 | value function |
| Total Impact | Yes | Yes | No |
| Marginal Impact | Yes | Yes | No |
| Mutual Impact | No | No | No |
| Minimum Impact | Yes | Yes | Yes |
| Differential Total Impact | Yes | Yes | No |

Table 7.11: Can different scoring methods detect the strong colluding agents?

## Chapter 8

## Conclusion

In this section the contributions of this thesis is summarized and several avenues for future work are introduced.

### 8.1 Contributions

In this thesis we introduced the first implemented and successful automatic collusion detection system. This method ranks pairs of agents in order of the chance that they are colluding. This is done through a novel data structure, called a collusion table, which captures the impact that each agent has on the utility of all agents in the game. We showed how collusion tables can be created using value functions which estimate the expected value of each game state for each player. Three different methods for constructing a value function are introduced and two of them were implemented. We showed how to use collusion tables to detect collusion by introducing five different scoring methods which rank pairs of players in order of the possibility that they are colluding given the collusion values.

Poker was utilized as our test domain because of its well-defined formal structure and high similarities to a real-world system in which collusion is prohibited. We developed a population of colluding and non-colluding agents in this domain and created a synthetic dataset of collusive behaviour to evaluate our technique. We evaluated our collusion detection method and showed that our method can successfully detect the colluding agents in our dataset.

This work has two main contributions. First, we developed an automated collu-
sion detection method and extended the applicability of the advantage sum idea to the domain of collusion detection. This technique does not require any hand-crafted features or depend on any characteristics of the domain. Second, we designed a general collusion method and constructed colluding agents that gain more utility through their collusion in our test domain. We created a synthetic dataset of agents using this method, which could be used in future research on collusion.

### 8.2 Future Work

There are three main directions for future work:

- We tested our technique using [2-4] Hold'em poker which is a smaller variation of Texas Hold'em poker. Therefore, it can be easily modified to evaluate corpus of games in Annual Computer Poker Competition in which bots are playing Texas Hold'em poker. This work was a first step to detecting collusion in human settings. The next step would be to evaluate this technique in a real-world corpus. Human datasets could be huge in size. A collusion detection system might need to employ a high level filtering of the data first, and then collusion tables could be constructed. Although finding sufficient labelled human data is a challenge, we hope to be able to find such data and apply this method in the near future.
- In this thesis, the PIVAT value function was introduced and the theoretical advantage of using it is shown. We are interested in implementing this technique and evaluating its effectiveness in future collusion detection methods.
- Finally, this method can be extended to other sequential decision making domains, since it does not rely on any characteristics of poker. It would be interesting to investigate reformulating our approach to be used in the regulation of market places in which detecting coalitions of parties is a real-world challenge. While there is a vast literature on preventing coalition by designing auctions in which collusion is not the preferable strategy, there have not been as many studies on detecting collusion when it might have occurred.

We are interested in extending our technique and similarly utilizing collusion tables and scoring methods to identify collusive behaviour in market places.

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## Appendix A

## Complete Results of Different Collusion Scores

A. 1 Experiment 1: Strong v.s. 5s Weak Agents

| Rank | Agent $i$ | Agent $j$ | Money |
| :--- | :--- | :--- | :---: |
| 63 | W5．CR | S．NC | -14.3355 |
| 64 | W5．CL | S．CR | -14.6155 |
| 65 | S．DF | W5．CR | -15.23325 |
| 66 | W5．DF | W5．CA | -15.98116667 |
| 67 | W5．CA | S．CR | -16.04820833 |
| 68 | S．PR | W5．CR | -16.14225 |
| 69 | W5．DF | W5．CB | -16.2335 |
| 70 | S．CR | W5．CB | -16.61233333 |
| 71 | W5．NC | W5．CA | -16.78558333 |
| 72 | W5．NC | W5．CB | -17.32033333 |
| 73 | W5．NC | W5．PR | -17.53195833 |
| 74 | W5．CA | W5．PR | -17.56933333 |
| 75 | W5．PR | W5．CB | -17.79466667 |
| 76 | W5．DF | W5．PR | -18.39758333 |
| 77 | W5．CR | S．CA | -18.846 |
| 78 | W5．CR | S．CB | -20.44754167 |
| 79 | W5．CL | S．CL | -20.86691667 |
| 80 | W5．CL | W5．DF | -22.57908333 |
| 81 | W5．CL | W5．NC | -23.06370833 |
| 82 | W5．CL | W5．PR | -23.479375 |
| 83 | W5．CL | W5．CB | -24.25791667 |
| 84 | W5．CL | W5．CA | -24.50579167 |
| 85 | W5．CR | S．CR | -28.86554167 |
| 86 | W5．CL | W5．CR | -29.55108333 |
| 87 | W5．CR | W5．PR | -30.32220833 |
| 88 | W5．DF | W5．CR | -36.97533333 |
| 89 | W5．NC | W5．CR | -37.356375 |
| 90 | W5．CA | W5．CR | -39.35104167 |
| 91 | W5．CR | W5．CB | -39.417625 |
|  |  |  |  |


|  | $\begin{aligned} & n \\ & \infty \\ & \infty \\ & \underset{\sim}{n} \\ & \infty \\ & \infty \end{aligned}$ | 8.523208333 |  | $\begin{aligned} & n \\ & 0 \\ & 0 \\ & 2 \\ & \vdots \\ & \infty \end{aligned}$ |  | $\begin{aligned} & n \\ & \text { N } \\ & 0 \\ & 0 \\ & \\ & \end{aligned}$ | $7.091333333$ |  | m m n N N $\infty$ $\infty$ n | $3.784120833$ |  |  |  | $\stackrel{n}{n}$ |  |  | $\begin{aligned} & n \\ & \underset{1}{n} \\ & \hdashline \\ & \infty \\ & 0 \\ & - \end{aligned}$ |  |  | $\begin{aligned} & n \\ & \underset{1}{n} \\ & 0 \\ & 0 \\ & 1 \\ & 1 \end{aligned}$ | $\begin{aligned} & m \\ & \text { m } \\ & \text { m } \\ & \infty \\ & \infty \\ & \infty \\ & 0 \\ & \dot{p} \end{aligned}$ |  | m m m m m on 子 n | $\begin{aligned} & \hat{6} \\ & \mathbf{o} \\ & \underset{\sim}{2} \\ & \underset{i}{2} \\ & \underset{i}{2} \end{aligned}$ | $\begin{aligned} & \text { N} \\ & \underset{\lambda}{2} \\ & \text { 人̀ } \end{aligned}$ |  |  | $\begin{aligned} & n \\ & n \\ & n \\ & n \\ & n \\ & n \\ & n \\ & n \end{aligned}$ | $$ |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | $\begin{aligned} & \frac{a}{2} \\ & \underset{\sim}{n} \\ & 3 \end{aligned}$ | $\begin{aligned} & n \\ & 0 \\ & n \\ & 3 \end{aligned}$ | $\begin{aligned} & u \\ & Z \\ & \dot{u} \end{aligned}$ | $\begin{aligned} & \mathbb{U} \\ & \dot{U} \end{aligned}$ | $\begin{aligned} & 0 \\ & \dot{3} \\ & 3 \end{aligned}$ | $\begin{aligned} & \frac{n}{n} \\ & n \\ & n \\ & 3 \end{aligned}$ | $\begin{aligned} & \frac{\alpha}{2} \\ & n \\ & n \\ & 3 \end{aligned}$ | $\begin{aligned} & \sim \\ & \sim \\ & \sim \end{aligned}$ | $\frac{\alpha}{\frac{\alpha}{n}}$ | $\underset{\sim}{\sim}$ | $\begin{aligned} & \frac{x}{2} \\ & 0 \\ & n \\ & 3 \end{aligned}$ | $\begin{gathered} \sim \\ \sim \\ \dot{N} \end{gathered}$ | $\begin{aligned} & \infty \\ & \dot{v} \\ & \dot{n} \end{aligned}$ | $\begin{aligned} & \mathbb{Z} \\ & u \\ & u \\ & i \\ & z \end{aligned}$ | $\begin{aligned} & \infty \\ & 0 \\ & i \\ & i \\ & i \end{aligned}$ | $\begin{aligned} & U \\ & Z \\ & \dot{v} \end{aligned}$ | $\stackrel{\sim}{\alpha}$ | $\begin{aligned} & \mathbb{U} \\ & \dot{N} \end{aligned}$ | $\begin{aligned} & \infty \\ & \dot{v} \\ & \dot{n} \end{aligned}$ | $\frac{\alpha}{\alpha}$ | $\begin{aligned} & \frac{1}{2} \\ & 2 \\ & 3 \\ & 3 \end{aligned}$ | $\begin{aligned} & \frac{2}{2} \\ & 2 \\ & n \\ & 3 \end{aligned}$ | $\begin{aligned} & U \\ & Z \\ & n \\ & B \\ & B \end{aligned}$ | $\begin{aligned} & \mathbb{Z} \\ & \underset{3}{2} \\ & B \end{aligned}$ | $\begin{aligned} & \underline{\sim} \\ & \underset{3}{n} \\ & \dot{3} \end{aligned}$ | $\begin{aligned} & U \\ & \text { n } \\ & 3 \end{aligned}$ | $\begin{aligned} & \mathscr{U} \\ & \dot{\sim} \end{aligned}$ | $\begin{aligned} & \sim \\ & \sim \\ & u \end{aligned}$ | $\begin{aligned} & \frac{\alpha}{2} \\ & n \\ & n \\ & 3 \end{aligned}$ | 1 3 2 |
|  | $\stackrel{\rightharpoonup}{\mathrm{O}}$ | $\underset{\sim}{\mathbb{N}}$ | $\begin{aligned} & \underset{U}{U} \\ & \dot{v} \end{aligned}$ | $\begin{aligned} & U \\ & 0 \\ & i n \\ & 3 \end{aligned}$ | $\dot{Q}$ | $\begin{aligned} & Q \\ & \dot{v} \\ & \dot{\sim} \end{aligned}$ | $\begin{gathered} \mathbb{U} \\ \dot{\omega} \end{gathered}$ | $\begin{aligned} & 1 \\ & 2 \\ & 10 \\ & 3 \end{aligned}$ | $\begin{aligned} & \cup \\ & Z \\ & \vdots \\ & B \\ & B \end{aligned}$ | $\frac{\alpha}{\alpha}$ | $\stackrel{\alpha}{2 \times 1}$ | $\begin{aligned} & \mathbb{U} \\ & \dot{U} \end{aligned}$ | $\begin{aligned} & \underset{U}{u} \\ & \dot{v} \end{aligned}$ | $\frac{\alpha}{2}$ | $\frac{\alpha}{n}$ |  | $\begin{aligned} & 3 \\ & 10 \\ & 3 \\ & 3 \end{aligned}$ | $\begin{aligned} & 0 \\ & n \\ & 3 \\ & 3 \end{aligned}$ | $\begin{aligned} & u \\ & u \\ & n \\ & 3 \end{aligned}$ | $\begin{aligned} & 3 \\ & 20 \\ & 3 \end{aligned}$ | $\begin{aligned} & U \\ & \dot{Q} \end{aligned}$ | $\begin{aligned} & U \\ & \dot{u} \end{aligned}$ | $\begin{aligned} & U \\ & \dot{v} \end{aligned}$ | ¢ |  | $\dot{N}$ | $\begin{aligned} & \cup \\ & Z \\ & \vdots \\ & \vdots \\ & B \end{aligned}$ | $\begin{aligned} & 1 \\ & 2 \\ & 10 \\ & 3 \end{aligned}$ | $\begin{aligned} & \underset{U}{\cup} \\ & \dot{\sim} \end{aligned}$ | ${ }^{1}$ |
| $\stackrel{\substack{\pi}}{\sim}$ | n | m | $\cdots$ | $\cdots$ | ल | $\cdots$ | ले | ¢ | $\ni$ | $\stackrel{\text { ¢ }}{ }$ | $\stackrel{\text { が }}{+}$ | 寸 | ワ | $\stackrel{\bigcirc}{+}$ | － | $\stackrel{\infty}{+}$ | ） | 안 | $\stackrel{\sim}{n}$ | N | n | － | $\cdots$ | $\stackrel{\circ}{\bullet}$ | in | $\stackrel{\infty}{\sim}$ | 合 | 8 | $\bigcirc$ | $\bigcirc$ |


| Rank | Agent $i$ | Agent $j$ | Money |
| :---: | :---: | :---: | :---: |
| 1 | S．CA | S．CB | 86.35875 |
| 2 | S．DF | S．CA | 39.023 |
| 3 | S．DF | S．NC | 38.92045833 |
| 4 | S．CA | S．NC | 37.22775 |
| 5 | S．DF | S．CB | 36.64808333 |
| 6 | S．NC | S．CB | 36.55991667 |
| 7 | S．PR | S．CB | 30.48295833 |
| 8 | S．PR | S．CA | 30.06370833 |
| 9 | S．PR | S．NC | 29.96825 |
| 10 | S．DF | S．PR | 29.54375 |
| 11 | W5．CA | W5．CB | 24.03591667 |
| 12 | S．CL | S．DF | 17.74708333 |
| 13 | S．CL | S．NC | 17.648625 |
| 14 | S．CL | S．CR | 17.43766667 |
| 15 | S．CL | S．CA | 15.15670833 |
| 16 | S．CL | S．CB | 14.83466667 |
| 17 | W5．NC | S．NC | 11.95775 |
| 18 | S．DF | W5．DF | 11.906375 |
| 19 | S．CL | S．PR | 11.8565 |
| 20 | S．DF | W5．NC | 11.709125 |
| 21 | W5．DF | S．CB | 11.52475 |
| 22 | S．DF | W5．CB | 11.12375 |
| 23 | W5．NC | S．CA | 11.053875 |
| 24 | W5．DF | S．NC | 11.02204167 |
| 25 | W5．NC | S．CB | 10.75641667 |
| 26 | W5．CA | S．NC | 10.53875 |
| 27 | W5．DF | S．CA | 10.34166667 |
| 28 | S．DF | W5．CA | 10.08333333 |
| 29 | S．NC | W5．CB | 9.952 |
| 30 | S．DF | S．CR | 9.400333333 |
| 31 | S．NC | W5．PR | 9.285958333 |

Table A．1：E1－Agent pairs ranked according to the total money gained（mbb／g）．


| $=\begin{gathered} t \\ \substack{0 \\ n \\ n \\ 0 \\ 0 \\ i \\ \hline} \end{gathered}$ | $\left\{\begin{array}{l} n \\ \vdots \\ \vdots \\ \vdots \\ \vdots \\ \vdots \\ i \end{array}\right.$ | 0 <br>  <br>  <br>  <br> 0 <br> 0 <br> 0 <br> 0 | $\begin{aligned} & \infty \\ & \infty \\ & + \\ & \vdots \\ & \infty \\ & 0 \\ & 0 \\ & 0 \end{aligned}$ | $\begin{gathered} 2 \\ 2 \\ \infty \\ 0 \\ 0 \\ 0 \\ i \end{gathered}$ |  |  | $\begin{aligned} & \text { n } \\ & \text { Ò } \\ & 0 \end{aligned}$ | $\begin{aligned} & \infty \\ & \infty \\ & \infty \\ & \underset{\infty}{\infty} \\ & \hat{c} \\ & \vdots \\ & i \end{aligned}$ | $\begin{gathered} 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \end{gathered}$ | $\begin{aligned} & n \\ & \text { a } \\ & 0 \\ & \vdots \\ & \vdots \\ & 0 \\ & i \end{aligned}$ | $\begin{aligned} & 0 \\ & 0 \\ & 0 \\ & 0 \\ & 0 \\ & 0 \\ & 0 \\ & i \end{aligned}$ | $\begin{aligned} & 2 \\ & \frac{2}{2} \\ & \frac{0}{2} \\ & 0 \\ & i \end{aligned}$ | $\begin{aligned} & n \\ & \underset{\sim}{2} \\ & \\ & 0 \\ & 0 \\ & i \end{aligned}$ | $$ | o 0 0 0 0 0 0 0 | $\left\{\begin{array}{l} \underset{n}{n} \\ \underset{\sim}{n} \\ \\ \\ i \end{array}\right.$ | $\infty$ $\infty$ $\vdots$ n 0 0 0 $i$ | $\begin{aligned} & 0 \\ & n \\ & \infty \\ & n \\ & 0 \\ & 0 \\ & 0 \end{aligned}$ | $\infty$ $n$ $n$ $n$ 0 0 0 $i$ | $\begin{aligned} & 2 \\ & 2 \\ & \\ & 0 \\ & 0 \\ & 0 \\ & 0 \end{aligned}$ | I 0 0 0 0 0 0 0 0 |  | $\begin{aligned} & n \\ & n \\ & n \\ & \infty \\ & 0 \\ & 0 \\ & i \end{aligned}$ | $\begin{aligned} & 0 \\ & 0 \\ & 0 \\ & 0 \\ & 0 \\ & 0 \\ & 0 \\ & i \end{aligned}$ |  | $\left\lvert\, \begin{gathered} \mathbb{N} \\ \mathbf{y} \\ 0 \\ 0 \\ 0 \\ 0 \\ 1 \end{gathered}\right.$ | $\begin{aligned} & 0 \\ & 0 \\ & 0 \\ & 0 \\ & 0 \\ & 0 \\ & 0 \end{aligned}$ |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | $\left\{\begin{array}{l} \frac{\alpha}{n} \\ n \\ n \\ 3 \end{array}\right.$ | $\begin{aligned} & U \\ & Z \\ & n \\ & 3 \\ & 3 \end{aligned}$ | $\begin{aligned} & 0 \\ & 0 \\ & n \\ & 0 \\ & 0 \end{aligned}$ |  |  | $$ | $\begin{aligned} & \mathrm{L} \\ & \underset{\sim}{n} \\ & \underset{3}{2} \end{aligned}$ | $\stackrel{\rightharpoonup}{\mathrm{a}}$ |  | $\begin{aligned} & U \\ & Z \\ & n \\ & 3 \end{aligned}$ | $\begin{aligned} & 0 \\ & Z \\ & 2 \\ & 3 \\ & Z \end{aligned}$ | $\begin{aligned} & 4 \\ & u \\ & i \\ & 3 \end{aligned}$ | $\begin{aligned} & 4 \\ & U \\ & \dot{n} \\ & 3 \end{aligned}$ | $\begin{aligned} & 0 \\ & 10 \\ & i n \\ & 3 \end{aligned}$ | $\left\{\begin{array}{l} n \\ \vdots \\ \vdots \\ n \\ n \end{array}\right.$ | $\begin{aligned} & U \\ & Z \\ & n \\ & n \\ & 3 \end{aligned}$ | $\stackrel{0}{0}$ | $\begin{aligned} & \cup \\ & Z \\ & \dot{\sim} \end{aligned}$ | $\begin{aligned} & \text { U } \\ & \text { Un } \\ & 1 \end{aligned}$ |  |  | $\begin{aligned} & 0 \\ & 0 \\ & n \\ & 3 \end{aligned}$ | $\begin{aligned} & m \\ & 0 \\ & n \\ & z \end{aligned}$ |  | $\begin{aligned} & 0 \\ & 3 \\ & 3 \end{aligned}$ | $\begin{aligned} & 2 \\ & n \\ & 3 \\ & 3 \end{aligned}$ |  |  |  |
|  | $\dot{v}$ | $\begin{gathered} 3 \\ \dot{c} \end{gathered}$ | $\begin{aligned} & 0 \\ & n \\ & 0 \\ & 0 \end{aligned}$ | $\underset{\sim}{\sim}$ | $\begin{aligned} & 0 \\ & 0 \\ & 3 \end{aligned}$ | $\begin{aligned} & \sim \\ & 0 \\ & \dot{\sim} \end{aligned}$ |  | $15$ | $3$ | $\dot{\omega}$ |  | $\stackrel{\rightharpoonup}{n}$ |  |  | $\dot{i} \dot{i}$ |  | $\begin{aligned} & \mathbb{U} \\ & \dot{n} \\ & \dot{3} \end{aligned}$ | $\begin{aligned} & 0 \\ & n \\ & n \\ & 3 \end{aligned}$ | $\vdots \begin{aligned} & 0 \\ & n \\ & 3 \end{aligned}$ | $\begin{aligned} & U \\ & Z \\ & \dot{W} \end{aligned}$ | $\begin{aligned} & \mathbb{U} \\ & \vdots \\ & n \\ & i \end{aligned}$ | $\underset{\sim}{\sim}$ | $\frac{\alpha}{2}$ | $\begin{aligned} & \mathbb{U} \\ & \dot{\sim} \\ & \dot{n} \end{aligned}$ |  | $\dot{v}$ | $\begin{gathered} \bar{i} \\ n \\ 3 \end{gathered}$ |  | $\begin{aligned} & 2 \\ & i \\ & \vdots \\ & 3 \\ & 3 \end{aligned}$ |
| $\stackrel{\substack{\tilde{n}\\}}{(N)}$ |  |  |  |  | N | $\infty$ | ¢ |  | F | $\stackrel{\text { }}{ }$ |  |  | $\stackrel{\sim}{4}$ | $\bigcirc$ | ¢ | $\stackrel{\infty}{+}$ | － | $\bigcirc$ | 的 | に |  | 守 |  |  |  |  |  |  | V |


| Rank | Agent $i$ | Agent $j$ | TI |
| :--- | :--- | :--- | :---: |
| 1 | S．CA | S．CB | 1.00555 |
| 2 | W5．CA | W5．CB | 0.385842 |
| 3 | S．CR | W5．CR | 0.375167 |
| 4 | S．CL | S．CR | 0.349336 |
| 5 | W5．CL | S．CR | 0.091698 |
| 6 | S．CL | W5．CR | 0.0797685 |
| 7 | S．CL | S．DF | 0.0432276 |
| 8 | S．CR | S．PR | 0.0418236 |
| 9 | S．CL | S．NC | 0.0393955 |
| 10 | S．CA | S．DF | 0.0376595 |
| 11 | S．CB | S．DF | 0.0374919 |
| 12 | S．CL | S．CB | 0.0366992 |
| 13 | S．CL | S．CA | 0.0331157 |
| 14 | S．CL | S．PR | 0.0308204 |
| 15 | S．CA | S．NC | 0.0284244 |
| 16 | S．CA | S．PR | 0.0283438 |
| 17 | W5．CL | S．CB | 0.0270478 |
| 18 | S．CB | S．PR | 0.0265555 |
| 19 | S．CB | S．NC | 0.0262797 |
| 20 | S．DF | S．NC | 0.0252142 |
| 21 | W5．CL | S．NC | 0.0227646 |
| 22 | W5．CL | S．DF | 0.0213757 |
| 23 | W5．CL | S．CA | 0.02031 |
| 24 | S．DF | S．PR | 0.00878357 |
| 25 | S．NC | S．PR | 0.0080518 |
| 26 | S．CR | W5．PR | 0.00407007 |
| 27 | S．CL | W5．DF | -0.00850434 |
| 28 | W5．CL | W5．DF | -0.00999859 |
| 29 | W5．CL | S．PR | -0.0114101 |
| 30 | W5．CL | W5．NC | -0.0128569 |
| 31 | W5．CL | W5．PR | -0.0152997 |
|  |  |  |  |
| 1 |  |  |  |

Table A．2：E1－Agent pairs ranked according to the Total Impact Score．


Table A.3: E1-Agent pairs ranked according to the Marginal Impact Score.


| Rank | Agent $i$ | Agent $j$ | MiI | Rank | Agent $i$ | Agent $j$ | MiI |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | S.CA | S.CB | 0.414292 | 32 | W5.DF | S.CB | -0.0582174 |
| 2 | S.CL | S.CR | 0.0964158 | 33 | S.CL | S.PR | -0.0585908 |
| 3 | W5.CA | W5.CB | 0.0406547 | 34 | W5.NC | S.CA | -0.0585909 |
| 4 | S.DF | S.NC | 0.00535241 | 35 | S.DF | S.CR | -0.0587645 |
| 5 | S.NC | S.CB | -0.00691819 | 36 | W5.NC | S.PR | -0.0590812 |
| 6 | S.DF | S.CA | -0.00702479 | 37 | S.CL | S.CB | -0.059765 |
| 7 | S.CA | S.NC | -0.0071533 | 38 | W5.NC | S.CB | -0.0599747 |
| 8 | S.DF | S.CB | -0.00797061 | 39 | S.CA | W5.CB | -0.0618488 |
| 9 | S.DF | S.PR | -0.0136976 | 40 | W5.CL | W5.CB | -0.0619993 |
| 10 | S.PR | S.NC | -0.0137525 | 41 | S.CR | S.NC | -0.0624509 |
| 11 | S.PR | S.CB | -0.0218562 | 42 | W5.DF | S.CA | -0.0636885 |
| 12 | S.PR | S.CA | -0.0223335 | 43 | S.PR | S.CR | -0.0637637 |
| 13 | S.CL | S.DF | -0.0413776 | 44 | S.CB | W5.CB | -0.0638164 |
| 14 | W5.CL | S.NC | -0.0442192 | 45 | S.PR | W5.PR | -0.0650275 |
| 15 | S.CL | S.NC | -0.0461937 | 46 | W5.CA | S.CA | -0.0650894 |
| 16 | W5.CL | S.CR | -0.0462927 | 47 | W5.CL | W5.NC | -0.06545 |
| 17 | W5.CL | S.DF | -0.0467339 | 48 | S.CL | S.CA | -0.0662283 |
| 18 | S.DF | W5.NC | -0.0479013 | 49 | W5.CL | W5.DF | -0.0677114 |
| 19 | W5.NC | S.NC | -0.0518318 | 50 | W5.CL | W5.CA | -0.0697317 |
| 20 | S.DF | W5.DF | -0.0525409 | 51 | W5.CL | W5.PR | -0.0708804 |
| 21 | W5.CL | S.CB | -0.0528611 | 52 | S.DF | W5.PR | -0.071911 |
| 22 | S.DF | W5.CA | -0.0532672 | 53 | W5.CL | S.PR | -0.0723838 |
| 23 | S.DF | W5.CB | -0.0536039 | 54 | S.NC | W5.PR | -0.0747727 |
| 24 | S.PR | W5.CA | -0.0541288 | 55 | S.CA | W5.PR | -0.0768113 |
| 25 | S.NC | W5.CB | -0.0546656 | 56 | S.CB | W5.PR | -0.0768821 |
| 26 | W5.DF | S.NC | -0.0560043 | 57 | W5.DF | S.CR | -0.0769805 |
| 27 | W5.CA | S.NC | -0.0565428 | 58 | S.CL | W5.DF | -0.0804949 |
| 28 | W5.CA | S.CB | -0.0566923 | 59 | W5.NC | S.CR | -0.082862 |
| 29 | W5.CL | S.CA | -0.0568909 | 60 | W5.DF | W5.CA | -0.0846052 |
| 30 | W5.DF | S.PR | -0.0571369 | 61 | W5.NC | W5.CA | -0.0869093 |
| 31 | S.PR | W5.CB | -0.0578862 | 62 | S.CL | W5.CA | -0.0871059 |

Table A.4: E1-Agent pairs ranked according to the Min Impact Score.

|  | 0 <br>  <br> 7 <br>  <br> 0 <br> 0 <br> 1 |  |  |  |  |  | $\left\{\begin{array}{l} n \\ \frac{n}{0} \\ \frac{n}{0} \\ 0 \end{array}\right.$ | $\begin{aligned} & N \\ & N \\ & N \\ & n \\ & \vdots \\ & i \end{aligned}$ | $\begin{aligned} & \underset{\sim}{2} \\ & \underset{N}{n} \\ & \underset{\sim}{i} \end{aligned}$ | $\pm$ 0 0 $n$ $\vdots$ $\vdots$ $i$ |  |  | $\begin{aligned} & 5 \\ & 5 \\ & 0 \\ & 0 \\ & 0 \\ & 0 \\ & 0 \\ & 0 \\ & 0 \\ & 0 \\ & \hline \end{aligned}$ | $\begin{aligned} & 2 \\ & 2 \\ & \frac{2}{2} \\ & \frac{2}{0} \end{aligned}$ | $\left\{\begin{array}{l} 2 \\ y \\ 7 \\ 0 \\ 0 \end{array}\right.$ | $\mathfrak{l}$ | $\begin{aligned} & \text { O} \\ & \text { O} \\ & \text { on } \\ & \text { N} \\ & \vdots \end{aligned}$ | $\begin{aligned} & 0 \\ & 0 \\ & 0 \\ & 0 \\ & \vdots \\ & \vdots \end{aligned}$ |  | $\begin{gathered} 0 \\ \substack{0 \\ \vdots \\ \vdots \\ i \\ \hline} \end{gathered}$ | $\begin{aligned} & \text { N } \\ & \text { N} \\ & \\ & \hline \end{aligned}$ |  |  |  |  | $\begin{aligned} & 0 \\ & 0 \\ & 0 \\ & 0 \\ & 0 \\ & 0 \\ & 0 \end{aligned}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | $$ |  |  | $\begin{array}{l\|l} \infty & \infty \\ 0 & 0 \\ n & 1 \\ 3 & 3 \\ \hline \end{array}$ | $\underset{\sim}{n}$ | $\begin{array}{r} \infty \\ 0 \\ n \\ 3 \\ 3 \end{array}$ | $\left\{\begin{array}{l} 2 \\ 20 \\ 3 \\ 3 \end{array}\right.$ | $\begin{aligned} & \mathbb{U} \\ & \dot{v} \end{aligned}$ |  |  |  | 0 0 $i n$ 3 3 |  | $\left\{\begin{array}{l} \frac{\alpha}{2} \\ \\ \end{array}\right.$ | $\begin{aligned} & \frac{2}{3} \\ & \text { n } \\ & 3 \end{aligned}$ | $\pi$ | $25$ |  | $\begin{aligned} & \text { č } \\ & n \\ & 3 \end{aligned}$ | $\xrightarrow{\sim}$ |  |  |  | U | $\xrightarrow{0}$ |
|  |  | $\begin{array}{\|c\|c} 1 \\ i n \\ i & 3 \\ i & 0 \end{array}$ | $$ |  | $\underset{\sim}{\widetilde{U}} \underset{\sim}{U}$ |  | $\left\{\begin{array}{l} n \\ \\ n \\ n \\ n \end{array}\right.$ | O | $\begin{aligned} & 4 \\ & u \\ & n \\ & 3 \\ & 3 \end{aligned}$ | $\begin{aligned} & \mathbb{U} \\ & n \\ & n \\ & 3 \\ & 1 \end{aligned}$ |  | $3$ | $\dot{3}$ | $18$ |  |  | $\begin{aligned} & 0 \\ & n \\ & 3 \end{aligned}$ | $\stackrel{\rightharpoonup}{i}$ |  |  | $\begin{aligned} & \mathbb{U} \\ & \dot{v} \end{aligned}$ | $\underset{\sim}{U}$ |  | $\stackrel{\sim}{3}$ | 告 | U |
|  |  | む ¢ |  | ¢ 0 | 08 |  | $\wedge$ |  |  |  | へ |  |  |  | $\infty$ |  | N | $\infty$ | － |  | $\infty$ | ふ |  | 2 | \％ | a |


|  |  | $\begin{aligned} & n \\ & \underset{\sim}{2} \\ & \underset{\sim}{\infty} \\ & 0 \\ & \hline \end{aligned}$ | $\begin{aligned} & 1 \\ & \substack{1 \\ 0 \\ 2 \\ 0 \\ 0 \\ 0 \\ 0 \\ \hline} \end{aligned}$ | $\begin{aligned} & 7 \\ & \sqrt{n} \\ & \stackrel{n}{2} \\ & 0 \\ & 0 \end{aligned}$ | $\begin{aligned} & n \\ & 0 \\ & 0 \\ & 0 \\ & 0 \\ & 0 \\ & 0 \\ & 0 \\ & i \end{aligned}$ | $\begin{aligned} & 0 \\ & 0 \\ & 0 \\ & 0 \\ & 0 \\ & 0 \end{aligned}$ | $\begin{aligned} & \infty \\ & 0 \\ & 0 \\ & 0 \\ & 0 . \\ & 0 \\ & i \end{aligned}$ | $\left\{\begin{array}{l} \frac{d}{2} \\ \underset{\lambda}{2} \\ \underset{o}{3} \end{array}\right.$ | $:$ |  | $\begin{aligned} & \text { I } \\ & \text { İ } \\ & \underset{\sim}{1} \end{aligned}$ | $\begin{aligned} & \underset{N}{7} \\ & \underset{\sim}{7} \\ & 0 \end{aligned}$ | $\begin{aligned} & \vec{N} \\ & \underset{\sim}{2} \\ & \underset{\sim}{1} \end{aligned}$ | $\begin{aligned} & n \\ & 2 \\ & 2 \\ & \underset{\sim}{n} \\ & 0 \end{aligned}$ | $\begin{aligned} & \infty \\ & \infty \\ & \underset{\sim}{\infty} \\ & \hline 0 \\ & \hline 1 \end{aligned}$ | $\stackrel{N}{N}$ | $\begin{aligned} & \vec{n} \\ & 0 \\ & \cdots \\ & \vdots \\ & 0 \end{aligned}$ | $\frac{2}{2}$ | $\begin{aligned} & \hat{2} \\ & \hat{2} \\ & \underset{1}{2} \end{aligned}$ | $\begin{gathered} n \\ \tilde{N} \\ \tilde{y} \\ \underset{0}{0} \end{gathered}$ | $\begin{aligned} & 2 \\ & 0 \\ & \text { N} \\ & \text { a } \\ & \underset{0}{0} \end{aligned}$ | $\begin{aligned} & \underset{7}{7} \\ & \underset{\sim}{2} \\ & \underset{\sim}{3} \end{aligned}$ | $\begin{aligned} & \bar{o} \\ & \underset{\sim}{2} \\ & \underset{0}{2} \end{aligned}$ |  |  | $\begin{gathered} \underset{\sim}{2} \\ \underset{\sim}{1} \\ \underset{\sim}{9} \end{gathered}$ | $\begin{aligned} & \overrightarrow{2} \\ & 0 \\ & \underset{1}{2} \\ & \underset{0}{0} \end{aligned}$ | $\begin{aligned} & \infty \\ & 0 \\ & 0_{1} \\ & 1 \\ & \vdots \\ & 0 \end{aligned}$ | $0$ | $\begin{array}{l\|l} 1 & n \\ 0 & 0 \\ 0 & 0 \\ n & n \\ 0 & 0 \\ 0 & 0 \end{array}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | $\begin{aligned} & \vec{E} \\ & \stackrel{y}{4} \\ & \underset{4}{U} \\ & \dot{v} \end{aligned}$ | $\underset{\sim}{\infty}$ | $\begin{aligned} & U \\ & Z \\ & \vdots \end{aligned}$ | $\frac{\alpha}{2}$ | $\begin{aligned} & 1 \\ & \underset{n}{n} \\ & n \\ & 3 \end{aligned}$ | $\begin{aligned} & U \\ & Z \\ & n \\ & Z \end{aligned}$ | $\underset{\sim}{\infty}$ | $\dot{U}$ | $\underset{\sim}{\infty}$ | $\begin{aligned} & \mathbb{U} \\ & \dot{U} \\ & \dot{L} \end{aligned}$ | $\begin{aligned} & U \\ & Z \\ & n \\ & Z \end{aligned}$ | $\begin{aligned} & \text { Ư } \\ & \dot{\sim} \end{aligned}$ | $\stackrel{\sim}{\mathcal{U}} \underset{\dot{\sim}}{ }$ | $\begin{aligned} & \frac{x}{n} \\ & n \\ & n \\ & 3 \end{aligned}$ | $\underset{\sim}{\sim}$ | $\begin{aligned} & \underset{U}{U} \\ & \dot{\sim} \end{aligned}$ | $\begin{aligned} & \underset{U}{U} \\ & \dot{\sim} \end{aligned}$ | $\underset{\sim}{Q}$ | $\begin{aligned} & \frac{2}{2} \\ & n \\ & n \\ & 3 \end{aligned}$ | $\begin{aligned} & 9 \\ & n \\ & 3 \end{aligned}$ |  | $\begin{aligned} & \frac{x}{n} \\ & n \\ & n \\ & 3 \end{aligned}$ | $\begin{aligned} & \mathbb{U} \\ & \dot{\sim} \end{aligned}$ | $\begin{aligned} & \mathscr{U} \\ & \dot{n} \\ & 3 \end{aligned}$ | $\begin{aligned} & U \\ & Z \\ & n \\ & 3 \end{aligned}$ | $\begin{aligned} & U \\ & Z \\ & \dot{N} \end{aligned}$ | $\begin{aligned} & \frac{x}{2} \\ & n \\ & n \\ & 3 \end{aligned}$ | $\begin{aligned} & \frac{x}{n} \\ & n \\ & 3 \\ & 3 \end{aligned}$ | $\begin{aligned} & \mathbb{U} \\ & \dot{N} \\ & \dot{3} \end{aligned}$ | $$ |
|  |  | $\stackrel{L}{\mathrm{~L}}$ | $\begin{aligned} & U \\ & \dot{n} \\ & 3 \end{aligned}$ | $\begin{aligned} & 0 \\ & n \\ & n \\ & \end{aligned}$ | $\begin{aligned} & \mathrm{O} \\ & \dot{x} \end{aligned}$ | $\begin{aligned} & \text { In } \\ & \underset{\sim}{n} \\ & n \end{aligned}$ | $\underset{\sim}{U}$ | $\begin{aligned} & u \\ & Z \\ & n \\ & 3 \end{aligned}$ | $: \begin{aligned} & u \\ & Z \\ & n \\ & 3 \end{aligned}$ | نِ | نِ | $\begin{aligned} & \vec{n} \\ & \vdots \\ & \end{aligned}$ | $\frac{\alpha}{\underset{\sim}{2}}$ | $\begin{aligned} & \underset{\sim}{U} \\ & \dot{\sim} \end{aligned}$ | $\begin{aligned} & 1 \\ & \\ & \\ & \\ & \hline \end{aligned}$ | $\begin{aligned} & U \\ & Z \\ & n \\ & n \\ & 3 \end{aligned}$ | $\stackrel{\text { L }}{\substack{a \\ ~}}$ | $\begin{aligned} & 0 \\ & \dot{n} \\ & \vdots \end{aligned}$ | $\begin{aligned} & U \\ & Z \\ & n \\ & 3 \end{aligned}$ | $\begin{aligned} & \text { U } \\ & n \\ & 3 \end{aligned}$ | $\begin{aligned} & 0 \\ & \dot{3} \\ & 3 \end{aligned}$ | $\begin{aligned} & 0 \\ & n \\ & 3 \end{aligned}$ | $\begin{aligned} & 9 \\ & n \\ & n \end{aligned}$ | $\dot{U}$ | $\begin{aligned} & 2 \\ & n \\ & i \\ & 3 \end{aligned}$ |  | $\dot{U}$ | $\begin{aligned} & \vec{a} \\ & n \\ & n \\ & 3 \end{aligned}$ | $3$ | $\begin{array}{ll} u \\ 0 & 0 \\ n \\ n & n \\ 3 \end{array}$ |
|  | $\stackrel{\pi}{\sim}$ | $\cdots$ | m | n | $\cdots$ | n | $\infty$ | 2 | $\bigcirc$ | F | ～ | $\stackrel{\square}{7}$ | $\pm$ | ヶ | $\bigcirc$ | － | $\stackrel{\infty}{+}$ | $\bigcirc$ | $\bigcirc$ | 云 | $\cdots$ | n | $\stackrel{+}{\square}$ | $n$ | $\bigcirc$ | in | $\infty$ | 눈 | O | 6 O |


| Rank | Agent i | Agent | DI |
| :--- | :--- | :--- | :---: |
| 1 | S．CA | S．CB | 0.59348 |
| 2 | W5．CR | S．CR | 0.281338 |
| 3 | W5．CA | W5．CB | 0.262688 |
| 4 | S．CL | S．CR | 0.0502188 |
| 5 | S．DF | S．PR | 0.00218301 |
| 6 | S．DF | S．NC | -0.00397064 |
| 7 | S．PR | S．NC | -0.00730324 |
| 8 | S．CL | S．DF | -0.012755 |
| 9 | S．CL | S．NC | -0.0238906 |
| 10 | W5．DF | S．PR | -0.0394925 |
| 11 | S．DF | W5．DF | -0.0442297 |
| 12 | W5．CL | S．CR | -0.0449244 |
| 13 | S．DF | W5．NC | -0.0454422 |
| 14 | S．CL | S．PR | -0.0463691 |
| 15 | W5．NC | S．PR | -0.0484037 |
| 16 | S．PR | W5．PR | -0.0513328 |
| 17 | W5．NC | S．NC | -0.0535115 |
| 18 | W5．DF | S．NC | -0.0563818 |
| 19 | W5．CL | S．DF | -0.0587852 |
| 20 | W5．CL | S．NC | -0.0607848 |
| 21 | S．CA | S．NC | -0.0674986 |
| 22 | S．DF | W5．PR | -0.0692615 |
| 23 | S．PR | W5．CA | -0.0695348 |
| 24 | S．DF | W5．CA | -0.0717318 |
| 25 | S．NC | S．CB | -0.0724677 |
| 26 | S．PR | W5．CB | -0.0740849 |
| 27 | S．NC | W5．PR | -0.0777367 |
| 28 | S．DF | W5．CB | -0.0785399 |
| 29 | S．NC | W5．CB | -0.0795029 |
| 30 | S．PR | S．CA | -0.0807079 |
| 31 | S．PR | S．CB | -0.0817253 |
|  |  |  |  |

Table A．5：E1－Agent pairs ranked according to the Differential Total Impact Score．

| $\frac{3}{2}$ |  | $\begin{aligned} & \text { J } \\ & \text { İ } \\ & \text { N } \\ & \text { O} \\ & 0 \end{aligned}$ | $\begin{aligned} & \hat{1} \\ & \infty \\ & 0 \\ & 0 \\ & 0 \\ & 0 \\ & 0 \\ & 0 \end{aligned}$ | $$ | $\left\lvert\, \begin{aligned} & 9 \\ & y \\ & 0 \\ & 0 \\ & 0 \\ & 0 \\ & 0 \end{aligned}\right.$ | $\begin{aligned} & 0 \\ & 0 \\ & 0 \\ & \text { O} \\ & 0 \\ & 0 \\ & 0 \end{aligned}$ | $\left\{\begin{array}{l} m \\ n \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \end{array}\right.$ | $\begin{aligned} & 1 \\ & 2 \\ & \infty \\ & 0 \\ & 1 \\ & 0 \\ & 0 \\ & 1 \end{aligned}$ | $\begin{aligned} & 0 \\ & 2 \\ & 2 \\ & i \\ & 0 \\ & 0 \end{aligned}$ | 0 0 0 0 0 0 1 |  | $\begin{aligned} & 2 \\ & 0 \\ & 2 \\ & \vdots \\ & 0 \\ & 0 \\ & 0 \end{aligned}$ | $\begin{aligned} & 1 \\ & \\ & \\ & \\ & \vdots \\ & 0 \\ & 0 \\ & i \end{aligned}$ |  | $\begin{aligned} & 0 \\ & 0 \\ & \\ & \\ & 0 \\ & 0 \\ & 0 \end{aligned}$ | $\begin{aligned} & m \\ & \infty \\ & 0 \\ & 0 \\ & 0 \\ & 0 \\ & i \end{aligned}$ | 2 2 $n$ 0 0 0 0 $i$ | $\mathfrak{c}$ |  | $\begin{aligned} & n \\ & n \\ & n \\ & \infty \\ & 0 \\ & 0 \\ & 0 \\ & 0 \end{aligned}$ | $\begin{aligned} & n \\ & \infty \\ & 0 \\ & 0 \\ & 0 \\ & 0 \\ & 0 \\ & 0 \end{aligned}$ | $\begin{aligned} & \bar{n} \\ & n \\ & n \\ & \infty \\ & 0 \\ & 0 \\ & i \end{aligned}$ | $\begin{aligned} & n \\ & \infty \\ & n \\ & \infty \\ & \infty \\ & 0 \\ & 0 \\ & 0 \end{aligned}$ | $\begin{aligned} & \text { n } \\ & \hat{2} \\ & 2 \\ & 8 \\ & 0 \\ & 0 \end{aligned}$ | $\begin{gathered} \sim \\ n \\ \infty \\ \cdots \\ \vdots \\ i \end{gathered}$ |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\left\|\begin{array}{c} \overrightarrow{0} \\ 00 \\ 80 \end{array}\right\|$ |  |  | $\begin{aligned} & u \\ & Z \\ & n \\ & i \end{aligned}$ | $\begin{aligned} & x \\ & \underset{3}{n} \\ & 3 \end{aligned}$ | $\left\{\begin{array}{l} \alpha \\ 2 \\ n \\ n \end{array}\right.$ | $\begin{aligned} & n \\ & n \\ & n \\ & n \end{aligned}$ | $\begin{aligned} & \underset{U}{\sim} \\ & n \\ & 3 \end{aligned}$ | $\begin{aligned} & 0 \\ & 0 \\ & n \\ & 3 \end{aligned}$ | $\begin{aligned} & U \\ & Z \\ & i \end{aligned}$ | $\begin{aligned} & \mathbb{U} \\ & \dot{n} \\ & 3 \end{aligned}$ | $\stackrel{\rightharpoonup}{\mathrm{L}}$ | $\stackrel{\sim}{0}$ | $\begin{aligned} & U \\ & Z \\ & Z \end{aligned}$ | $\dot{U}$ | $\frac{\alpha}{2}$ | $\dot{C}$ | $\underset{\sim}{\sim}$ | $\dot{y}$ | $\begin{aligned} & U \\ & Z \\ & \dot{B} \end{aligned}$ | $\begin{aligned} & U \\ & Z \\ & \dot{n} \end{aligned}$ | $\underset{\sim}{\underset{\sim}{U}}$ | $\frac{\alpha}{2 \times 1}$ | $\frac{\alpha}{2 \times 1}$ | $\frac{\alpha}{\underset{\sim}{2}}$ | $\stackrel{\rightharpoonup}{\mathrm{L}}$ | $\stackrel{\rightharpoonup}{\Delta}$ |  |  |
| $\left\|\begin{array}{c} \stackrel{3}{2} \\ 80 \\ 80 \end{array}\right\|$ |  |  | $\begin{aligned} & n \\ & n \\ & 3 \end{aligned}$ | $\Omega$ | $\begin{aligned} & 0 \\ & \dot{v} \end{aligned}$ | $3$ | $3$ | $3$ | $\stackrel{I}{1}$ | $\mathfrak{l}$ |  | $\left\{\begin{array}{l} 0 \\ \dot{n} \\ \dot{n} \\ i \end{array}\right.$ |  | $\left\{\begin{array}{l} B \\ \dot{n} \\ i \end{array}\right.$ | $\left\lvert\, \begin{aligned} & \text { Ch } \\ & \dot{C} \end{aligned}\right.$ |  | $\begin{aligned} & \mathbb{U} \\ & \dot{u} \end{aligned}$ | $\dot{i} \dot{\dot{i}}$ | $\begin{aligned} & \mathbb{U} \\ & \dot{U} \end{aligned}$ | $\underset{\sim}{\sim}$ | $\underset{\sim}{\sim}$ | $\begin{aligned} & \mathbb{U} \\ & \dot{\sim} \end{aligned}$ | $\begin{aligned} & \text { O} \\ & \underset{\sim}{\sim} \end{aligned}$ | $\underset{\sim}{\infty}$ | $\begin{aligned} & \mathbb{U} \\ & \dot{\sim} \end{aligned}$ | $\begin{aligned} & \sim \\ & 0 \\ & \dot{\sim} \end{aligned}$ |  |  |
|  | \% | 2 | $\bigcirc$ | T | ${ }_{\circ}^{\circ}$ | \% |  |  | N | $\cdots$ | さ | $\cdots$ | $\bigcirc$ |  | $\infty$ |  |  |  | $\infty$ | $\infty$ | $\pm$ | $\infty$ | $\infty$ | - | $\infty$ | $\infty$ |  | ন |



| Rank | Agent $i$ | Agent $j$ | MuI |
| :--- | :--- | :--- | :---: |
| 1 | S.CR | W5.CR | 0.8205 |
| 2 | W5.CR | W5.NC | 0.177952 |
| 3 | W5.CR | W5.DF | 0.17314 |
| 4 | W5.CA | W5.CR | 0.158265 |
| 5 | W5.CB | W5.CR | 0.148728 |
| 6 | W5.CR | S.DF | 0.145593 |
| 7 | W5.CR | S.NC | 0.14041 |
| 8 | S.PR | W5.PR | 0.0898325 |
| 9 | S.CA | S.CB | 0.0864108 |
| 10 | S.CB | W5.CR | 0.0828874 |
| 11 | W5.CR | W5.PR | 0.0819924 |
| 12 | S.DF | W5.PR | 0.0799518 |
| 13 | W5.NC | W5.PR | 0.0775214 |
| 14 | S.CA | W5.CR | 0.0747294 |
| 15 | S.NC | W5.PR | 0.0720325 |
| 16 | W5.DF | W5.PR | 0.0711563 |
| 17 | W5.CA | W5.CB | 0.0684351 |
| 18 | W5.DF | W5.NC | 0.0651412 |
| 19 | W5.CB | W5.PR | 0.0565469 |
| 20 | S.DF | W5.NC | 0.0531545 |
| 21 | W5.CA | W5.PR | 0.0516538 |
| 22 | W5.CB | W5.NC | 0.0514842 |
| 23 | S.DF | W5.DF | 0.0514578 |
| 24 | W5.DF | S.PR | 0.0489385 |
| 25 | S.NC | W5.NC | 0.0480998 |
| 26 | W5.CA | W5.NC | 0.0475637 |
| 27 | W5.CR | S.PR | 0.0455969 |
| 28 | W5.NC | S.PR | 0.0442014 |
| 29 | W5.DF | S.NC | 0.0441112 |
| 30 | S.CR | W5.DF | 0.0409191 |
| 31 | W5.CA | W5.DF | 0.0390632 |
|  |  |  |  |

Table A.6: E1-Agent pairs ranked according to the Mutual Impact Score).

## A. 2 Experiment 2: Strong v.s. 2s Weak Agents

| Rank | Agent $i$ | Agent $j$ | Money |
| :--- | :--- | :--- | :---: |
| 63 | W2.CA | S.CR | -22.625375 |
| 64 | W2.PR | W2.CL | -24.96266667 |
| 65 | W2.CR | S.PR | -30.24016667 |
| 66 | W2.PR | W2.CA | -32.38191667 |
| 67 | W2.PR | W2.NC | -32.39558333 |
| 68 | W2.CR | S.DF | -32.888625 |
| 69 | W2.PR | W2.DF | -32.90866667 |
| 70 | W2.CR | S.NC | -33.10620833 |
| 71 | W2.PR | W2.CB | -33.20579167 |
| 72 | S.CA | W2.CR | -34.10908333 |
| 73 | S.CB | W2.CR | -35.47104167 |
| 74 | W2.CL | S.CR | -38.69241667 |
| 75 | W2.DF | W2.NC | -39.5385 |
| 76 | W2.CB | W2.DF | -40.419125 |
| 77 | W2.CB | W2.NC | -40.57116667 |
| 78 | W2.CA | W2.DF | -41.35133333 |
| 79 | W2.CA | W2.NC | -41.8875 |
| 80 | S.CL | W2.CR | -45.83 |
| 81 | W2.CL | W2.NC | -47.715 |
| 82 | W2.CL | W2.DF | -48.18916667 |
| 83 | W2.CA | W2.CL | -48.41 |
| 84 | W2.CB | W2.CL | -49.06916667 |
| 85 | W2.CR | S.CR | -63.22791667 |
| 86 | W2.CR | W2.DF | -68.47 |
| 87 | W2.CR | W2.NC | -69.78083333 |
| 88 | W2.CB | W2.CR | -69.97666667 |
| 89 | W2.CL | W2.CR | -70.15625 |
| 90 | W2.CA | W2.CR | -71.36291667 |
| 91 | W2.PR | W2.CR | -79.99791667 |
|  |  |  |  |

\begin{tabular}{|c|c|c|c|c|c|c|c|c|c|c|c|c|c|c|c|c|c|c|c|c|c|c|c|}
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\begin{array}{ll}
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0 \\
0 & 0 \\
0 & 2 \\
0 & 0 \\
0 \\
0 \\
0
\end{array}
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\\
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\vdots \\
\end{gathered}
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& 0 \\
& 0 \\
& 0 \\
& \underset{\sim}{7} \\
& 0
\end{aligned}
$$ \& $n$

$\cdots$

$n$ \&  \&  \& \[
$$
\begin{aligned}
& \infty \\
& \infty \\
& \infty \\
& \infty \\
& n \\
& \infty \\
& 0 \\
& 0 \\
& n
\end{aligned}
$$

\] \&  \&  \&  \&  \&  \&  \& \[

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\begin{aligned}
& \infty \\
& \infty \\
& \infty \\
& 0 \\
& \underset{\sim}{n} \\
& \underset{\sim}{1}
\end{aligned}
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\underset{i}{ }
\end{gathered}
$$

\] \&  \&  \&  \& \& \&  \& \& \& \& \[

\underset{\sim}{\sim}
\] <br>

\hline $$
\mathbb{8}
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\mathfrak{c | c}

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$$ \& \& \& $\cdots$ \& \& テ \& ヲ \& \& \& \& - \& $\stackrel{\infty}{+}$ \& \& \& \& \& \& \& \& $\bigcirc$ in \& \& \& \& No <br>

\hline
\end{tabular}

| Rank | Agent $i$ | Agent $j$ | Money |
| :--- | :--- | :--- | :---: |
| 1 | S.CA | S.CB | 119.9245833 |
| 2 | S.DF | S.NC | 83.48125 |
| 3 | S.CA | S.DF | 83.07791667 |
| 4 | S.CB | S.DF | 82.66416667 |
| 5 | S.CA | S.NC | 81.79208333 |
| 6 | S.CB | S.NC | 81.51625 |
| 7 | S.DF | S.PR | 75.52083333 |
| 8 | S.NC | S.PR | 75.03625 |
| 9 | S.CB | S.PR | 74.87333333 |
| 10 | S.CA | S.PR | 73.79083333 |
| 11 | S.CL | S.NC | 60.94833333 |
| 12 | S.CL | S.DF | 60.8625 |
| 13 | S.CA | S.CL | 57.37875 |
| 14 | S.CB | S.CL | 56.86333333 |
| 15 | S.CL | S.PR | 54.49875 |
| 16 | S.CL | S.CR | 50.11583333 |
| 17 | S.CR | S.NC | 42.21458333 |
| 18 | S.CR | S.DF | 41.6465 |
| 19 | S.CR | S.PR | 39.91570833 |
| 20 | S.CA | S.CR | 38.36758333 |
| 21 | S.CB | S.CR | 35.87870833 |
| 22 | W2.PR | S.NC | 9.155291667 |
| 23 | S.CA | W2.NC | 8.943416667 |
| 24 | S.DF | W2.NC | 8.631583333 |
| 25 | W2.CB | S.NC | 8.456083333 |
| 26 | W2.CB | S.CA | 7.609125 |
| 27 | W2.DF | S.DF | 7.2585 |
| 28 | W2.CB | S.DF | 7.155833333 |
| 29 | W2.CA | S.CB | 7.132833333 |
| 30 | W2.CA | S.NC | 7.121666667 |
| 31 | W2.DF | S.NC | 7.035 |
|  |  |  |  |

Table A.7: E2-Agent pairs ranked according to the total money gained (mbb/g).

| Rank | Agent $i$ | Agent $j$ | TI |
| :--- | :--- | :--- | :---: |
| 63 | W2．CL | S．CL | -0.204396 |
| 64 | W2．CB | S．CL | -0.209616 |
| 65 | W2．CB | W2．CA | -0.209993 |
| 66 | S．CA | W2．CL | -0.214823 |
| 67 | S．CB | W2．DF | -0.216394 |
| 68 | S．CB | W2．CL | -0.217498 |
| 69 | W2．NC | S．PR | -0.225511 |
| 70 | S．DF | W2．NC | -0.22914 |
| 71 | W2．DF | S．PR | -0.232071 |
| 72 | W2．NC | S．NC | -0.232615 |
| 73 | W2．CA | S．PR | -0.235737 |
| 74 | W2．CB | S．DF | -0.236154 |
| 75 | W2．CB | S．PR | -0.237098 |
| 76 | W2．DF | S．NC | -0.237308 |
| 77 | W2．DF | S．DF | -0.238279 |
| 78 | W2．CB | S．NC | -0.242364 |
| 79 | W2．CA | S．DF | -0.244706 |
| 80 | W2．CA | S．NC | -0.245536 |
| 81 | W2．CL | S．PR | -0.247188 |
| 82 | W2．CL | S．NC | -0.252549 |
| 83 | W2．CL | S．DF | -0.259712 |
| 84 | W2．CL | W2．CR | -0.291146 |
| 85 | W2．PR | W2．CR | -0.296062 |
| 86 | S．CL | W2．CR | -0.33153 |
| 87 | W2．CR | S．PR | -0.432061 |
| 88 | S．CA | W2．CR | -0.46662 |
| 89 | S．CB | W2．CR | -0.468154 |
| 90 | W2．CR | S．NC | -0.497572 |
| 91 | W2．CR | S．DF | -0.499057 |
|  |  |  |  |
|  |  |  |  |
| 75 |  |  |  |


|  |  |  | $\begin{aligned} & \text { N} \\ & 0 \\ & 0 \\ & 0 \\ & 0 \end{aligned}$ | $\begin{aligned} & \infty \\ & \underset{\sim}{\infty} \\ & \underset{0}{7} \end{aligned}$ | $\begin{aligned} & \infty \\ & \mathbf{Q}_{1}^{2} \\ & \underset{0}{3} \end{aligned}$ | $\left\lvert\, \begin{gathered} \left.\begin{array}{c} n \\ N \\ \\ 0 \end{array} \right\rvert\, \end{gathered}\right.$ | $\left\lvert\, \begin{aligned} & 0 \\ & \infty \\ & \infty \\ & 0 \\ & 0 \\ & 0 \\ & i \end{aligned}\right.$ | $\begin{aligned} & \underset{\sim}{2} \\ & \underset{0}{0} \end{aligned}$ | $\frac{\underset{2}{2}}{\frac{2}{9}}$ | $\begin{aligned} & a \\ & i \\ & b \\ & d \\ & i \end{aligned}$ | $\frac{\underset{2}{7}}{\frac{2}{2}}$ | $\begin{aligned} & \sqrt{n} \\ & \text { àn } \\ & \underset{0}{0} \\ & i \end{aligned}$ |  |  | $\left.\begin{gathered} 2 \\ 0 \\ 2 \\ \vdots \\ \vdots \\ 0 \\ i \end{gathered} \right\rvert\,$ | $\begin{aligned} & 9 \\ & \underset{y}{2} \\ & \stackrel{y}{2} \\ & \underset{i}{2} \end{aligned}$ | $\begin{aligned} & \infty \\ & 0 \\ & 0 \\ & \frac{2}{\infty} \\ & \frac{1}{0} \\ & \hline \end{aligned}$ | $\frac{2}{\infty}$ |  | $\begin{aligned} & \infty \\ & \infty \\ & \infty \\ & \infty \\ & \underset{i}{0} \\ & \hline 1 \end{aligned}$ | $\begin{aligned} & n \\ & 0 \\ & 0 \\ & 0 . \\ & 0 . \\ & 0 \end{aligned}$ | $\begin{aligned} & 7 \\ & \frac{7}{2} \\ & \frac{2}{3} \\ & \hline \end{aligned}$ | $\begin{gathered} n \\ n \\ 2 \\ \vdots \\ i \end{gathered}$ | $\begin{aligned} & \infty \\ & \infty \\ & \infty \\ & \frac{a}{o} \\ & i \end{aligned}$ |  | $\begin{gathered} \underset{\sim}{2} \\ \underset{0}{0} \\ \underset{0}{3} \end{gathered}$ | $\frac{9}{\mathrm{q}}$ | $\begin{aligned} & \infty \\ & \infty \\ & \infty \\ & \infty \\ & \stackrel{\infty}{\infty} \\ & \hline- \end{aligned}$ | $\begin{aligned} & 2 \\ & 0 \\ & 0 \\ & 0 \\ & \text { N} \\ & \text { in } \end{aligned}$ |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  | $\xrightarrow{\sim}$ | $\stackrel{\sim}{U}$ | $\begin{aligned} & U \\ & Z \\ & \dot{Z} \\ & \mathbf{Z} \end{aligned}$ | $\begin{aligned} & \cup \\ & Z \\ & i \\ & Z \end{aligned}$ | $\left\lvert\, \begin{aligned} & \text { U } \\ & Z \\ & i \\ & X \end{aligned}\right.$ | $\begin{aligned} & \text { U } \\ & \text { X } \end{aligned}$ | $\underset{\sim}{\mathcal{O}}$ | $\begin{aligned} & \mathbb{U} \\ & \dot{v} \end{aligned}$ | $\begin{aligned} & 1 \\ & \text { U } \\ & \text { X } \end{aligned}$ | $\begin{aligned} & \vec{i} \\ & \dot{Z} \end{aligned}$ |  | $\begin{aligned} & \text { A} \\ & \text { i } \\ & i \end{aligned}$ | $\underset{\sim}{U}$ | $\begin{array}{\|} \frac{\alpha}{2} \\ n_{i} \end{array}$ | $\begin{aligned} & \vec{i} \\ & \dot{X} \end{aligned}$ | $\stackrel{T}{\mathrm{~A}}$ | $\begin{aligned} & \underset{Z}{z} \\ & \underset{3}{3} \end{aligned}$ | $\begin{aligned} & \mathrm{U} \\ & \mathrm{i} \\ & \mathbf{X} \end{aligned}$ | $\begin{aligned} & \mathrm{I} \\ & \mathrm{~B} \\ & \mathrm{i} \end{aligned}$ | $\begin{aligned} & \underset{\sim}{U} \\ & \text { i } \\ & \text { B } \end{aligned}$ | $\begin{aligned} & z \\ & \lambda \\ & 3 \end{aligned}$ | $\begin{aligned} & Z \\ & i \\ & \mathcal{Z} \end{aligned}$ | $\underset{\sim}{\infty}$ | $\begin{aligned} & \mathbb{U} \\ & \dot{U} \end{aligned}$ | $\begin{aligned} & U \\ & Z \\ & \dot{Z} \\ & X \end{aligned}$ |  |  | $\underset{\sim}{\infty} \underset{\sim}{\sim}$ |
|  |  | $\begin{gathered} 0 \\ \dot{N} \\ i \\ i \end{gathered}$ | $\begin{aligned} & \mathbb{U} \\ & \dot{v} \end{aligned}$ | $\begin{aligned} & \infty \\ & \dot{\sim} \\ & \dot{\sim} \end{aligned}$ | $\begin{aligned} & \mathrm{I} \\ & \mathrm{~A} \\ & \mathrm{i} \\ & \mathrm{~B} \end{aligned}$ | $\left\{\begin{array}{l} \infty \\ 0 \\ i \\ z \end{array}\right.$ | $\begin{aligned} & \mathbb{K} \\ & \text { U } \\ & \text { i } \end{aligned}$ | $\begin{aligned} & \mathbb{U} \\ & \text { i } \\ & B \end{aligned}$ | $\begin{aligned} & \frac{\alpha}{2} \\ & \underset{i}{3} \end{aligned}$ | $\begin{aligned} & \frac{\alpha}{2} \\ & \dot{1} \\ & 3 \end{aligned}$ | $\begin{aligned} & 0 \\ & 0 \\ & \text { i } \\ & \text { in } \end{aligned}$ | $\begin{aligned} & \mathbb{4} \\ & \mathbf{N} \\ & \dot{i} \end{aligned}$ |  | $\begin{gathered} 0 \\ 2 \\ 2 \\ 3 \\ 3 \end{gathered}$ | $\begin{aligned} & \text { x } \\ & 2 \\ & \lambda \\ & 3 \end{aligned}$ | $\begin{array}{\|c} \underset{\sim}{x} \\ \underset{\sim}{2} \\ 3 \end{array}$ | $\underset{\sim}{U}$ | $\begin{aligned} & \alpha \\ & \underset{\sim}{2} \\ & \text { in } \end{aligned}$ | $$ | $\begin{aligned} & \mathbb{U} \\ & \dot{N} \\ & \mathbb{X} \end{aligned}$ | $\begin{aligned} & \underset{\sim}{U} \\ & \text { i } \\ & \underset{3}{2} \end{aligned}$ | $\begin{aligned} & \infty \\ & 0 \\ & \text { i } \\ & 3 \end{aligned}$ | $\begin{aligned} & \mathbb{U} \\ & \dot{v} \end{aligned}$ | $\dot{\Delta}$ | $\begin{aligned} & n \\ & 0 \\ & \text { N } \end{aligned}$ | $\begin{aligned} & \infty \\ & \mathcal{O}_{1} \\ & \dot{Z} \end{aligned}$ | $\underset{\sim}{\sim}$ | $\begin{aligned} & \mathbb{U} \\ & \text { i } \\ & \text { i } \end{aligned}$ | $\dot{i}$ | $\begin{aligned} & \mathbb{U} \\ & U \\ & \dot{N} \\ & \dot{B} \\ & \dot{N} \end{aligned}$ |
|  | $\underset{m}{2}$ |  | ¢ | m | \％ | － | $\infty$ | ले | ¢ | 子 | $\underset{\sim}{\text { }}$ | ๆ | 寸 | $\stackrel{\text { ® }}{\text { ¢ }}$ | $\bigcirc$ | － | $\stackrel{\infty}{+}$ | $\bigcirc$ | in | n | N | n | i | 尔 | $\stackrel{\sim}{\circ}$ | ¢ | $\stackrel{\sim}{\circ}$ | in | 8 | $\checkmark \mathrm{J}$ |


| Rank | Agent $i$ | Agent $j$ | TI |
| :--- | :--- | :--- | :---: |
| 1 | S．CA | S．CB | 1.13789 |
| 2 | S．CL | S．CR | 0.415116 |
| 3 | W2．CR | S．CR | 0.335429 |
| 4 | W2．PR | W2．CL | 0.17274 |
| 5 | W2．CB | S．CR | 0.116838 |
| 6 | W2．CA | S．CR | 0.11506 |
| 7 | S．CR | W2．NC | 0.0615105 |
| 8 | S．CR | W2．DF | 0.0489052 |
| 9 | S．CR | S．PR | 0.0101103 |
| 10 | S．CL | S．PR | -0.00386081 |
| 11 | S．CL | S．NC | -0.00430348 |
| 12 | W2．PR | S．CR | -0.00497141 |
| 13 | S．CA | S．CL | -0.00574327 |
| 14 | S．CB | S．PR | -0.0079995 |
| 15 | S．CL | S．DF | -0.00860491 |
| 16 | S．CA | S．PR | -0.0115699 |
| 17 | S．CB | S．CL | -0.0122188 |
| 18 | S．CB | S．DF | -0.025552 |
| 19 | S．DF | S．PR | -0.0258311 |
| 20 | S．NC | S．PR | -0.026532 |
| 21 | S．CA | S．DF | -0.0271918 |
| 22 | S．CB | S．NC | -0.0364653 |
| 23 | W2．PR | W2．NC | -0.03724 |
| 24 | S．DF | S．NC | -0.0415185 |
| 25 | S．CA | S．NC | -0.0427431 |
| 26 | W2．PR | W2．DF | -0.0443041 |
| 27 | W2．CL | S．CR | -0.06092 |
| 28 | W2．PR | W2．CB | -0.0709103 |
| 29 | W2．PR | W2．CA | -0.0785074 |
| 30 | S．CR | S．NC | -0.0824857 |
| 31 | S．CR | S．DF | -0.0831059 |

Table A．8：E2－Agent pairs ranked according to the Total Impact Score．

|  |  |  |  |  |  |  | $\begin{gathered} = \\ \vdots \\ \vdots \\ \vdots \\ \hline \end{gathered}$ |  |  | $0$ |  |  | $9$ | $\begin{gathered} a \\ \substack{a \\ 0 \\ 0 \\ 0 \\ \\ \vdots \\ \vdots \\ \vdots \\ i} \end{gathered}$ | $\hat{m}_{1}^{n}$ |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  | Cry | 菏 | $2$ |  | $2$ |  |  |  | >p | $p \mid$ |  | $\begin{aligned} & 0 \\ & i \\ & i \\ & i \\ & i \end{aligned}$ | $3$ | $18$ |  |  |  |  |  | - |
|  | B | U |  | $\begin{array}{lll} 0 & \\ i \\ i & 0 \\ 3 & 0 \\ i \end{array}$ | $\underset{\sim}{4} \underset{\sim}{\sim}$ |  |  |  | in | $\begin{aligned} & \mathbb{U} \\ & \dot{U} \end{aligned}$ |  |  |  | $5 \pi$ | $\begin{gathered} 0 \\ i \\ i \\ i \\ i \end{gathered}$ |  |  |  |  |  |  |  |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |


|  |  |  |  |  |  |  |  | $\underset{\sim}{\mathrm{F}}$ | $\stackrel{n}{n}$ | $\begin{aligned} & 2 \\ & 2 \\ & 2 \\ & 2 \\ & 2 \\ & 0 \\ & 0 \\ & 0 \\ & 0 \end{aligned}$ | $\underset{\sim}{m}$ |  |  | $1$ |  |  |  |  |  |  |  | $: \begin{gathered} \infty \\ \nabla \\ \hdashline \\ \hdashline \end{gathered}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | 込 |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  | r |  |  |  |  |  |  |  | 6 W |


| Rank | Agent $i$ | Agent $j$ | MaI |
| :--- | :--- | :--- | :---: |
| 1 | W2.CR | S.CR | 1.447774 |
| 2 | S.CA | S.CB | 1.16105 |
| 3 | S.CL | S.CR | 0.294364 |
| 4 | W2.CA | S.CR | 0.237159 |
| 5 | W2.CB | S.CR | 0.234338 |
| 6 | S.CR | W2.NC | 0.186118 |
| 7 | W2.PR | S.NC | 0.177526 |
| 8 | W2.CL | S.DF | 0.17574 |
| 9 | S.CR | W2.DF | 0.173203 |
| 10 | W2.CL | S.NC | 0.17293 |
| 11 | W2.CR | S.NC | 0.157832 |
| 12 | W2.CR | S.DF | 0.157477 |
| 13 | W2.PR | S.DF | 0.157206 |
| 14 | W2.CB | S.NC | 0.143451 |
| 15 | W2.CB | S.DF | 0.142371 |
| 16 | W2.DF | S.NC | 0.140003 |
| 17 | W2.CA | S.DF | 0.138915 |
| 18 | W2.NC | S.NC | 0.138124 |
| 19 | W2.CA | S.NC | 0.137902 |
| 20 | S.DF | W2.NC | 0.137902 |
| 21 | S.CL | W2.CR | 0.128606 |
| 22 | W2.DF | S.DF | 0.128345 |
| 23 | S.CA | W2.CL | 0.128026 |
| 24 | W2.PR | S.CA | 0.122675 |
| 25 | W2.PR | S.CB | 0.119151 |
| 26 | S.CB | W2.CL | 0.11602 |
| 27 | W2.CL | S.PR | 0.113557 |
| 28 | S.CA | W2.CR | 0.106443 |
| 29 | W2.CB | S.CB | 0.103871 |
| 30 | S.CA | W2.DF | 0.103496 |
| 31 | W2.CB | S.CA | 0.100685 |

Table A.9: E2-Agent pairs ranked according to the Marginal Impact Score.

| Rank | Agent $i$ | Agent $j$ | MiI | Rank | Agent $i$ | Agent $j$ | MiI | Rank | Agent $i$ | Agent $j$ | MiI |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | S.CA | S.CB | 0.457624 | 32 | W2.CA | W2.CL | -0.176971 | 63 | W2.DF | S.PR | -0.263105 |
| 2 | S.CL | S.CR | 0.11473 | 33 | W2.CB | W2.NC | -0.192993 | 64 | W2.CL | S.CL | -0.26402 |
| 3 | S.DF | S.NC | -0.0282926 | 34 | W2.DF | W2.NC | -0.195046 | 65 | S.CB | W2.DF | -0.265018 |
| 4 | W2.PR | W2.CL | -0.0345984 | 35 | W2.CB | W2.CL | -0.197297 | 66 | W2.CA | S.PR | -0.265206 |
| 5 | S.CB | S.DF | -0.0350939 | 36 | S.CB | S.CR | -0.198949 | 67 | W2.CB | S.PR | -0.271078 |
| 6 | S.DF | S.PR | -0.0355195 | 37 | S.CA | S.CR | -0.199039 | 68 | W2.CA | S.NC | -0.274076 |
| 7 | S.NC | S.PR | -0.0366603 | 38 | W2.CA | W2.DF | -0.199565 | 69 | S.DF | W2.NC | -0.274946 |
| 8 | S.CA | S.DF | -0.0389293 | 39 | W2.CB | W2.DF | -0.20283 | 70 | W2.CB | S.NC | -0.276407 |
| 9 | S.CB | S.NC | -0.0392255 | 40 | W2.CA | W2.NC | -0.203297 | 71 | W2.DF | S.DF | -0.276419 |
| 10 | S.CA | S.NC | -0.0394283 | 41 | W2.PR | S.CB | -0.205945 | 72 | W2.CB | S.DF | -0.276648 |
| 11 | S.CA | S.PR | -0.044461 | 42 | W2.PR | S.CA | -0.213517 | 73 | W2.NC | S.NC | -0.278451 |
| 12 | S.CL | S.PR | -0.0447222 | 43 | W2.PR | S.PR | -0.215025 | 74 | W2.DF | S.NC | -0.27997 |
| 13 | S.CB | S.PR | -0.0453031 | 44 | W2.PR | S.CL | -0.216585 | 75 | W2.CA | S.DF | -0.281203 |
| 14 | S.CL | S.DF | -0.0483226 | 45 | W2.PR | S.CR | -0.219509 | 76 | W2.CL | S.NC | -0.286641 |
| 15 | S.CL | S.NC | -0.0501218 | 46 | W2.PR | S.NC | -0.223101 | 77 | W2.CL | S.DF | -0.29082 |
| 16 | S.CA | S.CL | -0.0639879 | 47 | W2.CL | S.CR | -0.233194 | 78 | W2.CL | S.PR | -0.308512 |
| 17 | S.CB | S.CL | -0.0649373 | 48 | W2.PR | S.DF | -0.236853 | 79 | W2.CL | W2.CR | -0.361477 |
| 18 | W2.PR | W2.NC | -0.0855723 | 49 | W2.CB | S.CA | -0.242854 | 80 | W2.CA | W2.CR | -0.362905 |
| 19 | W2.PR | W2.DF | -0.0920056 | 50 | W2.CB | S.CB | -0.244613 | 81 | W2.CR | W2.NC | -0.369855 |
| 20 | W2.PR | W2.CB | -0.0963644 | 51 | W2.CA | S.CA | -0.246436 | 82 | W2.CR | W2.DF | -0.370178 |
| 21 | S.CR | S.PR | -0.102053 | 52 | W2.CA | S.CB | -0.247948 | 83 | W2.CB | W2.CR | -0.372931 |
| 22 | W2.PR | W2.CA | -0.102564 | 53 | S.CA | W2.NC | -0.248076 | 84 | W2.CB | W2.CA | -0.402951 |
| 23 | W2.CR | S.CR | -0.112098 | 54 | S.CB | W2.NC | -0.250251 | 85 | W2.PR | W2.CR | -0.452538 |
| 24 | S.CR | S.NC | -0.114776 | 55 | S.CL | W2.DF | -0.251583 | 86 | W2.CR | S.NC | -0.50848 |
| 25 | S.CR | S.DF | -0.11834 | 56 | S.CA | W2.DF | -0.254087 | 87 | W2.CR | S.DF | -0.511762 |
| 26 | S.CR | W2.NC | -0.142262 | 57 | W2.NC | S.PR | -0.254115 | 88 | S.CA | W2.CR | -0.531765 |
| 27 | S.CR | W2.DF | -0.14503 | 58 | W2.CA | S.CL | -0.254146 | 89 | S.CB | W2.CR | -0.532186 |
| 28 | W2.CL | W2.DF | -0.160158 | 59 | S.CL | W2.NC | -0.257257 | 90 | W2.CR | S.PR | -0.539223 |
| 29 | W2.CL | W2.NC | -0.160373 | 60 | S.CB | W2.CL | -0.257341 | 91 | S.CL | W2.CR | -0.580931 |
| 30 | W2.CB | S.CR | -0.162595 | 61 | S.CA | W2.CL | -0.25949 |  |  |  |  |
| 31 | W2.CA | S.CR | -0.164664 | 62 | W2.CB | S.CL | -0.262026 |  |  |  |  |

Table A.10: E2-Agent pairs ranked according to the Minimum Impact Score.

|  |  | $\begin{array}{ccc} - & 0 \\ \underset{\sim}{n} & 0 \\ N & \underset{n}{n} \\ & 0 \\ i & 0 \end{array}$ | $\begin{aligned} & \text { t} \\ & 0 \\ & 0 \\ & 0 \\ & 0 \\ & 0 \\ & 0 \end{aligned}$ |  | $\left[\begin{array}{l} 1 \\ 0 \\ 0 \\ 7 \\ 0 \\ 0 \end{array}\right.$ |  |  |  |  |  | $\begin{aligned} & 0 \\ & 0 \\ & 0 \\ & \text { n } \\ & \\ & \end{aligned}$ | $\begin{aligned} & n \\ & \underset{\sim}{n} \\ & \\ & \\ & \hline \end{aligned}$ |  |  | $\begin{gathered} 7 \\ \vdots \\ n \\ \substack{n \\ i \\ i} \end{gathered}$ | + 0 0 0 $\vdots$ $\vdots$ $i$ 1 |  | $\begin{aligned} & \infty \\ & \infty \\ & \infty \\ & \infty \\ & \vdots \\ & 1 \end{aligned}$ | $\begin{gathered} \hat{N} \\ \hat{N} \\ \underset{\sim}{\tau} \end{gathered}$ | $\begin{aligned} & \hat{0} \\ & 0 \\ & 0 \\ & 0 \\ & 0 \\ & 0 \\ & i \end{aligned}$ | $\begin{aligned} & n \\ & n \\ & n \\ & n \\ & \vdots \end{aligned}$ |  | $\begin{aligned} & 0 \\ & \substack{2 \\ \sim \\ ? \\ \hline} \end{aligned}$ |  |  | $\begin{aligned} & 1 \\ & 0 \\ & 0 \\ & 2 \\ & \\ & 0 \\ & i \end{aligned}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | $\begin{aligned} & 0 \\ & \text { in } \\ & \text { n } \\ & 2 \end{aligned}$ | $\begin{array}{ll} \alpha & - \\ \underset{\sim}{U} & \dot{X} \\ \dot{\alpha} & 3 \end{array}$ | $\begin{aligned} & Z \\ & \mathbf{Z} \\ & \mathbf{Z} \end{aligned}$ |  | $\begin{aligned} & \underset{i}{i} \\ & \underset{i}{i} \\ & 3 \end{aligned}$ |  |  |  | $\begin{aligned} & \mathbb{U} \\ & i \\ & X \\ & 3 \end{aligned}$ | $\begin{aligned} & Z \\ & \text { Z } \\ & \text { B } \end{aligned}$ | $\begin{aligned} & Z \\ & \text { i } \\ & z \end{aligned}$ | $\begin{aligned} & Z \\ & i \\ & z \end{aligned}$ | $\begin{array}{\|c} \mathrm{A} \\ \underset{X}{\mathrm{~N}} \\ \hline \end{array}$ | $\begin{aligned} & 0 \\ & \text { X } \\ & \text { B } \end{aligned}$ | $\begin{aligned} & \mathrm{O} \\ & \text { i } \\ & \text { B } \end{aligned}$ | $\begin{aligned} & u \\ & Z \\ & i \end{aligned}$ | $\left.\begin{gathered} 1 \\ s \\ i \end{gathered} \right\rvert\,$ | $\begin{aligned} & \mathrm{O} \\ & \text { i } \end{aligned}$ | Z | $\stackrel{\mathrm{i}}{\mathrm{i}}$ | $\begin{aligned} & \underset{U}{U} \\ & \text { X } \\ & \text { X } \end{aligned}$ | 2 |  |  |  |  |
| $\begin{array}{\|c} \stackrel{\rightharpoonup}{0} \\ \stackrel{0}{0} \\ \ll \end{array}$ |  |  | $\begin{aligned} & \vec{A} \\ & \stackrel{y}{2} \\ & \end{aligned}$ | $\begin{array}{lll} 1 & \infty \\ \text { L } & 0 \\ i & i \\ i & 3 \end{array}$ | $\begin{aligned} & \mathbb{U} \\ & 2 \\ & 2 \\ & z \end{aligned}$ | $$ | $\begin{aligned} & U \\ & \text { i } \\ & i \end{aligned}$ |  | $\begin{aligned} & \text { x } \\ & \\ & \end{aligned}$ | $\begin{aligned} & \text { O } \\ & \text { i } \\ & z \end{aligned}$ | $\begin{aligned} & \text { x } \\ & \\ & \text { n } \end{aligned}$ |  | $\begin{gathered} \underset{\sim}{2} \\ \underset{\sim}{3} \\ \end{gathered}$ | $\begin{aligned} & \mathbb{U} \\ & \text { i } \\ & \text { B } \end{aligned}$ | $\begin{aligned} & \mathrm{i} \\ & \mathrm{i} \end{aligned}$ | $2$ | $\begin{aligned} & \text { U } \\ & \text { i } \\ & \mathcal{Z} \end{aligned}$ | $\underset{\sim}{U}$ | $3$ | $\stackrel{\sim}{\mathrm{c}}$ | $\ll$ | N | U |  |  |  |
|  |  | す 勺 |  | ¢ |  | R | NN |  | + | $\cdots$ |  |  | $\stackrel{\sim}{\sim}$ |  |  |  |  | $\infty$ | $\pm$ | - | $\infty$ | $\stackrel{\infty}{\infty}$ | $\infty$ | இ |  | Ј |



| Rank | Agent $i$ | Agent $j$ | DI |
| :--- | :--- | :--- | :---: |
| 1 | S.CA | S.CB | 0.747346 |
| 2 | W2.CR | S.CR | 0.239798 |
| 3 | S.CL | S.CR | 0.170694 |
| 4 | S.CL | S.PR | 0.138459 |
| 5 | S.NC | S.PR | 0.129413 |
| 6 | S.DF | S.PR | 0.127802 |
| 7 | S.CL | S.DF | 0.114495 |
| 8 | S.CL | S.NC | 0.114352 |
| 9 | S.DF | S.NC | 0.0835895 |
| 10 | S.CA | S.PR | 0.0581982 |
| 11 | S.CB | S.PR | 0.0505685 |
| 12 | S.CA | S.CL | 0.0364487 |
| 13 | S.CB | S.CL | 0.023174 |
| 14 | S.CB | S.NC | 0.0152373 |
| 15 | S.CA | S.NC | 0.0079695 |
| 16 | S.CB | S.DF | 0.003627 |
| 17 | S.CA | S.DF | -0.000930967 |
| 18 | S.CR | S.PR | -0.0409137 |
| 19 | S.CR | S.DF | -0.064396 |
| 20 | S.CR | S.NC | -0.064556 |
| 21 | W2.CB | S.CR | -0.122357 |
| 22 | S.CR | W2.NC | -0.124974 |
| 23 | S.CR | W2.DF | -0.130104 |
| 24 | W2.CA | S.CR | -0.131269 |
| 25 | W2.PR | S.NC | -0.14029 |
| 26 | S.DF | W2.NC | -0.145776 |
| 27 | W2.DF | S.NC | -0.149715 |
| 28 | W2.CB | S.DF | -0.150351 |
| 29 | W2.CA | S.DF | -0.150617 |
| 30 | W2.NC | S.NC | -0.151966 |
| 31 | W2.CB | S.NC | -0.152511 |

Table A.11: E2-Agent pairs ranked according to the Differential Total Impact Score.

|  | $\begin{array}{r} 20 \\ 0 \\ 0 \\ 0 \\ 0 \end{array}$ | $39$ |  | $0^{9}$ | $99$ | $99$ | $0$ | $0$ |  |  | $0$ | pi |  |  |  | ${ }^{1}$ |  |  | $\top$ |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  |  |  |  |  |  |  |  |  |  |  | $\left(\begin{array}{c} n \\ n \\ n \end{array}\right)$ | $\frac{x}{2}$ |  |  |  |  |  |  |  |  |  | - |
|  |  | $\dot{r i v}$ |  |  |  |  |  |  |  | $\underset{c}{n} \begin{gathered} z \\ i \\ i \\ z \end{gathered}$ | O. |  | N. |  |  |  |  |  |  |  |  |  |  |  |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |



| Rank | Agent $i$ | Agent $j$ | MuI |
| :--- | :--- | :--- | :---: |
| 1 | W2.CR | S.CR | 1.1123 |
| 2 | W2.CR | S.DF | 0.656533 |
| 3 | W2.CR | S.NC | 0.655403 |
| 4 | S.CA | W2.CR | 0.57306 |
| 5 | S.CB | W2.CR | 0.558672 |
| 6 | W2.CR | S.PR | 0.525854 |
| 7 | S.CL | W2.CR | 0.460135 |
| 8 | W2.CL | S.DF | 0.435551 |
| 9 | W2.CL | S.NC | 0.42548 |
| 10 | W2.CB | S.NC | 0.385815 |
| 11 | W2.CA | S.DF | 0.38362 |
| 12 | W2.CA | S.NC | 00.383439 |
| 13 | W2.CB | S.DF | 0.378524 |
| 14 | W2.DF | S.NC | 0.377312 |
| 15 | W2.NC | S.NC | 0.370739 |
| 16 | S.DF | W2.NC | 0.367043 |
| 17 | W2.DF | S.DF | 0.366624 |
| 18 | W2.CL | S.PR | 0.360743 |
| 19 | S.CA | W2.CL | 0.34285 |
| 20 | W2.PR | S.NC | 0.340332 |
| 21 | W2.PR | S.DF | 0.339136 |
| 22 | S.CB | W.PL | 0.333516 |
| 23 | W2.DF | S.PR | 0.311935 |
| 24 | W2.CB | S.PR | 0.31141 |
| 25 | W2.CA | S.PR | 0.308956 |
| 26 | S.CA | W2.DF | 0.305581 |
| 27 | W2.NC | S.PR | 0.304444 |
| 28 | W2.CA | S.CA | 0.299093 |
| 29 | W2.CB | S.CB | 0.29885 |
| 30 | W2.CA | S.CB | 0.297202 |
| 31 | W2.CB | S.CA | 0.29692 |


[^0]:    ${ }^{1} h a$ refers to a history consisting of action $a$ concatenated to history $h$.

[^1]:    ${ }^{2}$ Agents used in the experiments of this thesis are developed by Computer Poker Research Group in University of Alberta (CPRG). The author's contribution is to modify them to collude and develop different kinds of players and agents.

[^2]:    ${ }^{1}$ http://www.computerpokercompetition.org/

